

D-FINITENESS: A SUCCESS STORY



Manuel Kauers · Institute for Algebra · JKU



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Definition.

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- A number α is called **algebraic** if there are integers c_0, \dots, c_d , not all zero, such that

$$c_0 + c_1\alpha + \dots + c_d\alpha^d = 0.$$

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- A function f is called **algebraic** if there are polynomials c_0, \dots, c_d , not all zero, such that

$$c_0(x) + c_1(x)f(x) + \dots + c_d(x)f(x)^d = 0.$$

Definition.

- A function f is called **D-finite** if there are polynomials p_0, \dots, p_r , not all zero, such that

$$p_0(x)f(x) + p_1(x)f'(x) + \dots + p_r(x)f^{(r)}(x) = 0.$$

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- A sequence (a_n) is called **D-finite** if there are polynomials q_0, \dots, q_s , not all zero, such that

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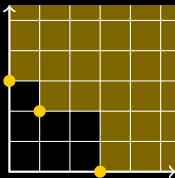
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$$(1-4x)f(x)^2 - 1 = 0$$

$$f(0) = 1$$

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 $1.1 < \alpha < 1.2$
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 $f(0) = 1$
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 $a_0 = 1$
- x^n $f_{n+1}(x) - x f_n(x) = 0$
 $xf'_n(x) - n f_n(x) = 0$
 $f_0(x) = 1$

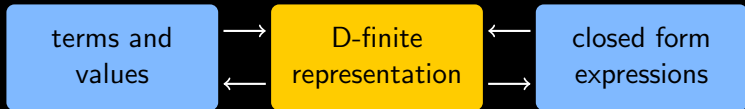
A computational ecosystem

D-finite
representation

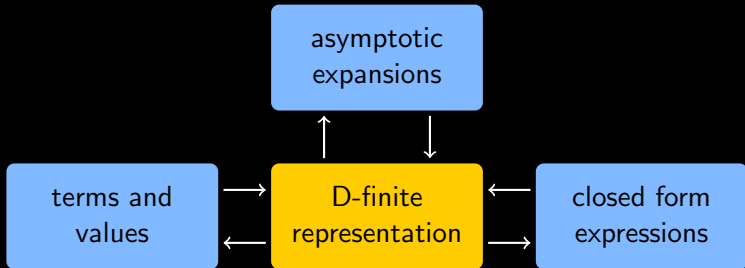
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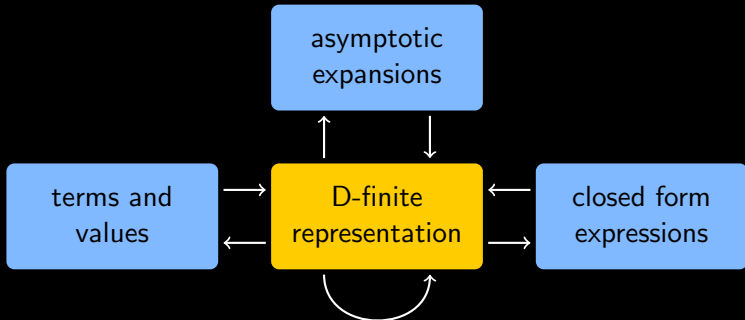
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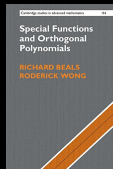
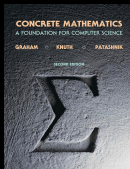
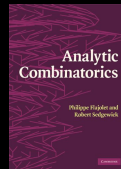
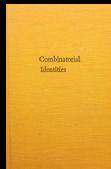
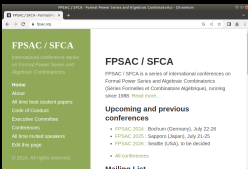
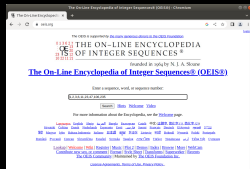
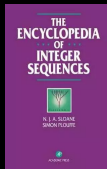
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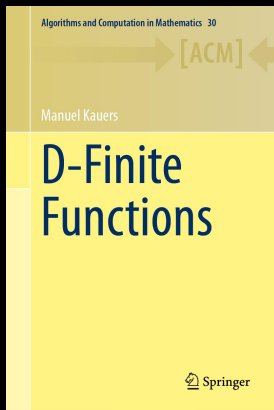


Who cares?



- Enumeration of lattice walks
- Permutation patterns
- Determinant evaluations
- Graph counting
- Analysis of algorithms
- Program verification
- Statistical mechanics
- Particle physics
- Special functions
- Analytic number theory
- Arithmetic number theory
- Experimental mathematics
- Numerical engineering
- Probability theory
- Knot theory
- Computational algebra
- Biology
- Coding theory
- Control theory
- Cryptography
- Statistics
- Spaceflight
- Sociology
- Simulation

What can we do?



Introduction

- Functions, sequences, and series
- D-Finiteness
- Applications
- Computer Algebra
- Guessing
- Hermite-Pade Approximation

The Recurrence Case in One Variable

- Evaluation
- The Solution Space
- Closure Properties
- Generalized Series Solutions
- Polynomial and Rational Solutions
- Hypergeometric and D'Alembertian Solutions

The Differential Case in One Variable

- Evaluation
- The Solution Space
- Closure Properties
- Generalized Series Solutions
- Polynomial and Rational Solutions
- Hyperexponential and D'Alembertian Solutions

Operators

- Ore Algebras and Ore Actions
- Common right divisors and left multiples
- Several functions
- Factorization
- Several variables
- Gröbner bases

Summation and Integration

- The indefinite problem
- The definite problem
- Further closure properties
- Creative telescoping
- Bounds
- Reduction-based algorithms



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$$D_x f(x, y) = \sum_{n, k \in \mathbb{Z}} k a_{n, k} x^{k-1} t^n$$

\ddots	\vdots	\vdots	\vdots	\ddots
\dots	$a_{-1,2}$	$a_{0,2}$	$a_{1,2}$	\dots
\dots	$a_{-1,1}$	$a_{0,1}$	$a_{1,1}$	\dots
\dots	$a_{-1,0}$	$a_{0,0}$	$a_{1,0}$	\dots
\dots	$a_{-1,-1}$	$a_{0,-1}$	$a_{1,-1}$	\dots
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\ddots	\vdots	\vdots	\vdots	\ddots

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In particular, $[x^{-1}]f$ is D-finite.

There are tons of ISSAC papers on how to compute such P, Q .

Do we really need this?

Example 1

Computer Algebra in the Service of Enumerative Combinatorics

Alin Bostan
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ABSTRACT

Counting lattice walks in restricted lattices is an important problem in enumerative combinatorics. Recently, computer algebra has been used to explore and to solve a number of difficult questions related to lattice walks. We give an overview of recent results on structural properties (e.g., algebraicity versus transcendence) and on explicit formulas for generating functions of walks with small steps in the quarter plane. In doing so, we emphasize the algorithmic nature of the methodology, especially two important paradigms “guess-and-prove” and “creative telescoping”.

CCS CONCEPTS

• Computing methodologies → Algebraic algorithms.

KEYWORDS

Computer algebra; Experimental mathematics; Guess-and-Prove; Creative telescoping; Enumerative combinatorics; Lattice paths; Generating functions; D-finite functions; Algebraic functions.

ACM Reference Format:

Alin Bostan, 2023. Computer Algebra in the Service of Enumerative Combinatorics. In *Proceedings of the 2023 International Symposium on Symbolic and Algebraic Computation (SAC’23)*, July 8–23, 2023, Virtual Event, Dublin, Ireland. ACM, New York, NY, USA, 8 pages. <https://doi.org/10.1145/3622433.3663207>

1 GENERAL PRESENTATION

1.1 Prelude

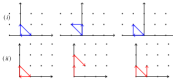
Consider the following innocent-looking problem.

A *random walk* is a path in \mathbb{Z}^2 taking steps from $\{T, \leftarrow, \rightarrow, \downarrow\}$ only. Show that, for any integer $n \geq 0$, the following quantities are equal:

(i) the number n_0 of random walks of length n in \mathbb{Z}^2 , ending at steps confined to the upper half plane $\mathbb{Z} \times \mathbb{N}$, then start and end at $(0, 0)$;

(ii) the number b_n of random walks of length n confined to the quarter plane \mathbb{N}^2 , that start at $(0, 0)$ and end on the diagonal $\mathbb{N} \times \mathbb{N}$.

For instance, when $n = 3$, this common value is $n_0 = b_3 = 3$, as shown in the next picture, obtained by exhaustive enumeration.



It appears that this problem is far from being trivial. Several solutions exist, but none of them is elementary. One of the main aims of the present text is to convince the reader that this problem (and many others with a similar flavor) can be solved with the help of a computer. More precisely, computer algebra tools can be used to discover and to prove the following equalities

$$a_{2n} = b_n = \frac{(3n)!}{n!^3} \quad \text{and} \quad a_n = b_n = 0 \quad \text{if } 3 \nmid n. \quad (1)$$

It goes without saying that such a simple and beautiful expression cannot be the result of mere chance. It turns out that closed form are quite rare for this kind of enumeration problem. Nevertheless, even in the absence of nice formulas, the structural properties of the corresponding enumeration sequences reflect the symmetries of the step set and of the evolution domain. Equation (1) shows that the sequences (a_n) and (b_n) are P-recursive, that is, they satisfy a linear recurrence with polynomial coefficients (in the index n). Equivalently, their generating functions $\sum_{n \geq 0} a_n x^n$ and $\sum_{n \geq 0} b_n x^n$ are D-finite, that is, they satisfy linear differential equations with polynomial coefficients (in the variable x). On the methodological (i.e., computer algebraic) side, one of the main messages that will emerge from the text is that, in the absence of closed formulas, the recurrence-differential equations themselves constitute the appropriate data structure to represent and manipulate P-recursive sequences and D-finite functions [40]. On the application (i.e., combinatorial) side, the main message is that these important properties of the enumeration sequences are intimately connected to the existence of a certain group, naturally attached to the step set $\{T, \leftarrow, \rightarrow, \downarrow\}$.

Example 2



How Does the Gerrymander Sequence Continue?

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Rutgers University (New Brunswick)
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Piscataway, NJ 08854-8019
USA
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Abstract

We compute a few additional terms of the gerrymander sequence (GDS sequence A340402) and provide closed equations for the generating functions of some sequences in the family.

Computational challenge for you (and your computer)

1 message

Doron Zeilberger <doronzeil@gmail.com>

Wed, May 4, 2022 at 3:43 PM

To: Manuel Kauers <manuel@kauers.de>

Cc: George Spahn <gs828@math.rutgers.edu>, Neil Sloane <njasloane@gmail.com>

Dear Manuel,

Hope you, Martina, and epsilon are doing well!

Can you (and your computers) meet the following challenge, in the secret url:<https://sites.math.rutgers.edu/~zeilberg/EM22/C27.pdf>

<https://sites.math.rutgers.edu/~zeilberg/ChessChallenge.txt>

If you do, I pledge to donate \$100 to the OEIS in your honor.

Also, if you can do it systematically, this may lead to a joint paper with my student who can do other boards.

Best wishes,
Doron

Gerrymandering (2), cont.

$T(k, d)$ = no. of ways to dissect a $k \times k$ square board
into d rook-connected regions of size k^2 / d .

(A348452, A348456,
A172477, A004003)

$k \backslash d$	1	2	3	4	5	6	7	8	9	...
1	1									
2	1	2	0	1						
3	1	0	10^a	0	0	0	0	0	1	
4	1	70^b	0	117	0	0	0	36^c	0	... 1@16
5	1	0	0	0	4096	0	0	0	0	... 1@25
6	1	80518	264500	442791	0	451206	0	0	✓	... 1@36
7	1	0	0	0	0	0	✓	0	0	... 1@49
8	1	?	0	?	0	0	0	✓	0	... 1@64

Most wanted: $T(8,2)$ = no. of ways to cut chessboard into 2 rook-connected regions of area 32

Ignore colors of chessboard squares; rotations, reflections count as different; regions need not have same shape.

How large will $T(8,2)$ be, roughly? How would you program it? How would you parallelize it?

Paul Zimmermann et al. in 2020 solved one of the RSA Challenge Problem,
It took them 2700 core years. How does $T(8,2)$ compare?








Gerrymandering (2), cont. $T(4,2) = 70$:







23 of 37
Math640.04.2022
34.6%

Gerrymandering (2), cont.

T(4,2) = 70:

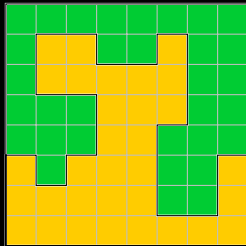
70 WAYS TO DISSECT 4x4 BOARD

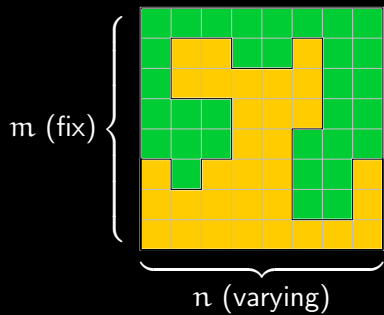
	CODE BEGINNING COUNT	#	
	2 2 2 2 B=4 C=0 G=4	# = 2	1
	1 1 3 3 B=6 C=2 G=2	# = 4	2
	1 2 2 3 B=6 C=4 G=2	# = 4	3
	1 3 1 3 B=10 C=6 G=2	# = 4	4
	1 2 1 4 B=8 C=5 G=1	# = 8	5
	1 3 2 2 B=7 C=4 G=1	# = 8	6
	1 2 3 2 B=7 C=6 G=1	# = 8	7

	1 1 1 4 B=7 C=3 G=1	# = 8	8
	1 3 3 1 B=8 C=4 G=2	# = 4	9
	1 1 2 4 B=6 C=3 G=2	# = 4	10
	1 1 3 2 B=9 C=6 G=1	# = 8	11
	2 1 3 2 B=8 C=6 G=2	# = 4	12
	1 1 1 3 B=10 C=4 G=2	# = 4	13
TOTAL = 70			

Gerrymandering (2), cont.

Tiling a Square with





Want:

$a_n =$

number of ways to split an $m \times n$ board
into two connected regions
of exactly the same size

Want: $a(t) = \sum_{n=0}^{\infty} a_n t^n$ where

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Detour:

$a_{n,k} =$ number of ways to split an $m \times n$ board
into (at most) two connected regions,
with exactly k yellow cells

Want: $a(t) = \sum_{n=0}^{\infty} a_n t^n$ where

$$a_n = \boxed{\begin{array}{c} \text{number of ways to split an } m \times n \text{ board} \\ \text{into two connected regions} \\ \text{of exactly the same size} \end{array}} = a_{n,nm/2}$$

Detour:

$$a_{n,k} = \boxed{\begin{array}{c} \text{number of ways to split an } m \times n \text{ board} \\ \text{into (at most) two connected regions,} \\ \text{with exactly } k \text{ yellow cells} \end{array}}$$

Want: $a(t) = \sum_{n=0}^{\infty} a_n t^n$ where

$$a_n = \boxed{\begin{array}{c} \text{number of ways to split an } m \times n \text{ board} \\ \text{into two connected regions} \\ \text{of exactly the same size} \end{array}} = a_{n,nm/2}$$

Detour: $a(x, t) = \sum_{n=0}^{\infty} \sum_{k=0}^{\infty} a_{n,k} x^k t^n$ where

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Key observation:

$$a(t) = [x^{-1}] \frac{1}{x} a(x^2, \frac{t}{x^m})$$

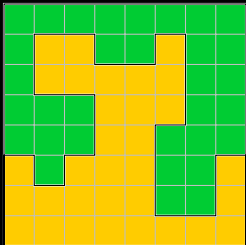
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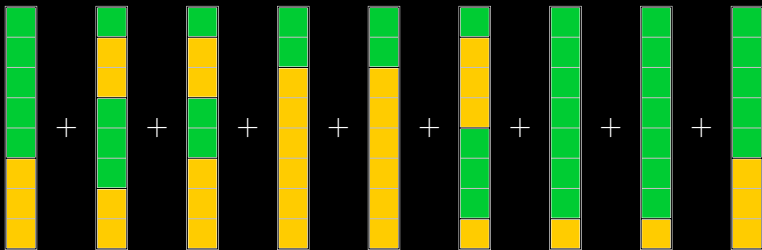
$$\begin{aligned} a(t) &= [x^{-1}] \underbrace{\frac{1}{x} a\left(x^2, \frac{t}{x^m}\right)} \\ &= \sum_{n=0}^{\infty} \sum_{k=0}^{\infty} a_{n,k} x^{2k-mn-1} t^n \end{aligned}$$

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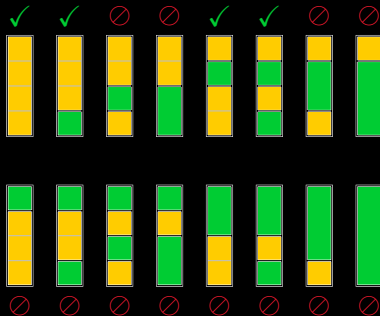
So we are done if we can find $a(x, t)$.

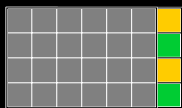




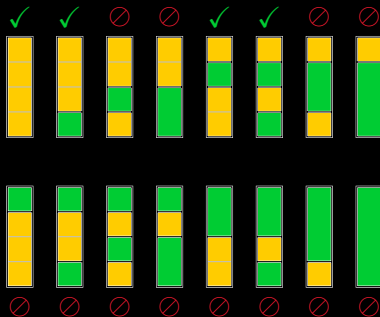


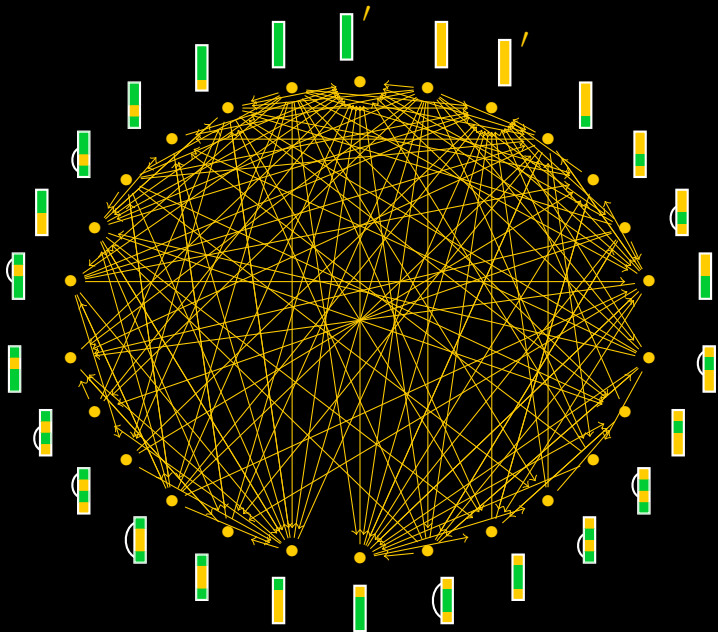
+ ?





+ ?





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[illegible]

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Challenge: Find a differential equation for $a(x)$ for some $m \geq 5$.



The number of ways to dissect an $m \times n$ board into two connected regions of equal size:

	1	2	3	4	5	6	7	8
1	0	1	0	1	0	1	0	1
2	1	2	3	4	5	6	7	8
3	0	3	0	19	0	85	0	355
4	1	4	19	70	245	856	2967	10164
5	0	5	0	245	0	8171	0	277969
6	1	6	85	856	8171	80518	806423	8059419
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Guttmann and Jensen showed that the diagonal is **not** D-finite.

Homework: Let

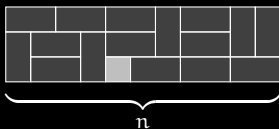
$a_n =$

number of tilings of a $3 \times n$ board
with dimers and exactly one monomer

Homework: Let

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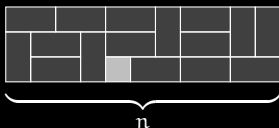
number of tilings of a $3 \times n$ board
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Homework: Let

$a_n =$

number of tilings of a $3 \times n$ board
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Show that (a_n) is D-finite and find a recurrence.

Example 3

Hardinian Arrays

Robert Dougherty-Bliss* Mannel Kauers[♯]

Submitted: Sep 1, 2023; Accepted: Mar 19, 2024; Published: Apr 5, 2024

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Abstract

In 2014, R.H. Hardin contributed a family of sequences about king-moves on an array to the On-Line Encyclopedia of Integer Sequences (OEIS). The sequences were recently noticed in an automated search of the OEIS by Kauers and Koutschan, who conjectured a recurrence for one of them. We prove their conjecture as well as some other conjectures stated in the OEIS entries. We also have some new conjectures for the asymptotics of Hardin's sequences.

Mathematics Subject Classifications: 05A15, 31F10

1 Introduction

The On-Line Encyclopedia of Integer sequences [15] contains over 350,000 sequences and perhaps tens of thousands of conjectures about them. Here we resolve some of these conjectures related to a family of sequences due to R.H. Hardin.

For any positive integer r , let $H_r(n, k)$ be the number of $n \times k$ arrays which obey the following rules:

- The entry in position $(1, 1)$ is 0, and the entry in position (n, k) is $\max(n, k) - r + 1$.
- The entry in position (i, j) must equal or be one more than each of the entries in positions $(i - 1, j)$, $(i, j - 1)$, and $(i - 1, j - 1)$.
- The entry in position (i, j) must be within r of $\max(i, j) - 1$.

We call such an arrangement of numbers a *Hardinian array*.

Equivalently, Hardinian arrays can be defined in terms of the king-distance between entries, i.e., the length of the shortest path that a king on a chessboard can take to get from one entry to the other. Using this notion, we can say that a Hardinian array has a fixed maximum king-distance between any two entries.

A253217 - OEIS

oeis.org/A253217

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0 1 3 6 2 7
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 : 20
 23 OE
 12 IS
 10 22 11 21

THE ON-LINE ENCYCLOPEDIA OF INTEGER SEQUENCES[®]

founded in 1964 by N. J. A. Sloane

[Hints](#)

(Greetings from [The On-Line Encyclopedia of Integer Sequences!](#))

A253217	<p>Number of $n \times n$ nonnegative integer arrays with upper left 0 and lower right its king-move distance away minus 2 and every value within 2 of its king move distance from the upper left and every value increasing by 0 or 1 with every step right, diagonally se or down.</p> <p>0, 0, 1, 19, 268, 3568, 47698, 649712, 9023385, 127419681, 1823918697, 26398702645, 385582981615, 5674890516295, 84060883775765, 1252066289632643, 18738613233957420, 281620474177057788, 4248088188086420832 (list; graph; refs; listen; history; text; internal format)</p> <p>OFFSET 1,4</p> <p>COMMENTS Diagonal of A253223.</p> <p>LINKS R. H. Hardin, Table of n, a(n) for n = 1..37 M. Kauers and C. Koutschan, Guessing with Little Data, ISSAC '22: Proceedings of the 2022 International Symposium on Symbolic and Algebraic Computation, July 2022, Pages 83-90. M. Kauers and C. Koutschan, Some D-finite and some possibly D-finite sequences in the OEIS, arXiv:2303.02793 [cs.SC], 2023. R. Dougherty-Bliss and M. Kauers, Hardinian Arrays, arXiv:2309.00487 [math.co], 2023.</p> <p>FORMULA Recurrence: $32*(1+n)*(1+2*n)^2*(161046+465785*n+551943*n^2+)$</p>
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Number of $n \times n$ nonnegative integer arrays with upper left 0 and lower right its king-move distance away minus 2 and every value within 2 of its king move distance from the upper left and every value increasing by 0 or 1 with every step right, diagonally SE or down.

$H_2(4,4) = 19$, because:

0001	0001	0001	0001	0001
0001	0001	0001	0001	0011
0001	,	0011	,	0111
1111	,	1111	,	1111

0001	0001	0001	0001	0001
0011	0011	0111	0111	1111
0111	,	1111	,	0111
1111	,	1111	,	1111

0011	0011	0011	0011	0011
0011	0011	0011	0111	0111
0011	,	0111	,	1111
1111	,	1111	,	1111

0011	0111	0111	0111
1111	0111	0111	1111
1111	,	0111	,
1111	,	1111	,

$H_2(n, n)$ begins as follows:

0, 0, 1, 19, 268, 3568, 47698, 649712, 9023385, 127419681, ...

$$\begin{aligned}
& (201600n^9 + 4942080n^8 + 53078112n^7 + 327661728n^6 + \\
& 1280700480n^5 + 3285342016n^4 + 5528828352n^3 + 5883447104n^2 + \\
& 3591093120n + 957662208)a(n) + (-970200n^9 - 24199560n^8 - \\
& 264810744n^7 - 1667830872n^6 - 6659340648n^5 - 17470825688n^4 - \\
& 30096410912n^3 - 32804461872n^2 - 20514211488n - \\
& 5603970816)a(n+1) + (589050n^9 + 14827590n^8 + 163756656n^7 + \\
& 1040895564n^6 + 4194035058n^5 + 11101344742n^4 + 19289250308n^3 + \\
& 21198776056n^2 + 13360158000n + 3676219776)a(n+2) + (294525n^9 + \\
& 7319295n^8 + 79828578n^7 + 501335472n^6 + 1997003589n^5 + \\
& 5229549731n^4 + 8997110634n^3 + 9799013608n^2 + 6125859120n + \\
& 1673566848)a(n+3) + (-121275n^9 - 3053295n^8 - 33716268n^7 - \\
& 214212552n^6 - 862421763n^5 - 2280190003n^4 - 3956305720n^3 - \\
& 4340670060n^2 - 2730542400n - 749859264)a(n+4) + (6300n^9 + \\
& 163890n^8 + 1863666n^7 + 12150660n^6 + 50023284n^5 + 134779202n^4 + \\
& 237527338n^3 + 263895164n^2 + 167643648n + 46381248)a(n+5) \stackrel{?}{=} 0
\end{aligned}$$

Let's prove this.

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But first, let's consider the case $r = 1$.

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Or maybe let's even start with $r = 0$.

0	1	2	3	4	5	6
1	1	2	3	4	5	6
2	2	2	3	4	5	6
3	3	3	3	4	5	6
4	4	4	4	4	5	6
5	5	5	5	5	5	6
6	6	6	6	6	6	6

0	1	2	3	4	5	6
1	1	2	3	4	5	6
2	2	2	3	4	5	6
3	3	3	3	4	5	6
4	4	4	4	4	5	6
5	5	5	5	5	5	6
6	6	6	6	6	6	6

$H_0(n, k) = 1$ for all n, k .

0

5


$$H_1(n, n) = ?$$

							5
							5
							5
							5
							5
							5
							5
5	5	5	5	5	5	5	5

$$H_1(n, n) = ?$$

0	1	2	2	3	4	5
1						5
2						5
3						5
3						5
4						5
5	5	5	5	5	5	5

$$H_1(n, n) = ?$$

							
	0	1	2	2	3	4	5
	1						5
	2						5
{	3						5
	3						5
	4						5
	5	5	5	5	5	5	5

$$H_1(n, n) = ?$$

							
0	1	2	2	3	4	5	
1	1	2	3	3	4	5	
2	2	2	3	3	4	5	
{	3	3	3	3	4	4	5
	3	3	4	4	4	5	5
4	4	4	4	5	5	5	
5	5	5	5	5	5	5	

$$H_1(n, n) = ?$$

0	1	2 2		3	4	5
1	1	2	3	3	4	5
2	2	2	3	3	4	5
3	3	3	3	4	4	5
3	3	4	4	4	5	5
4	4	4	4	5	5	5
5	5	5	5	5	5	5

$$H_1(n, n) = ?$$

$H_1(n, n)$ is the number of non-crossing lattice path tuples (P_1, \dots, P_{n-1}) where each path P_i starts somewhere on the left and ends somewhere at the top.

$H_r(n, n)$ is the number of non-crossing lattice path tuples (P_1, \dots, P_{n-r}) where each path P_i starts somewhere on the left and ends somewhere at the top.

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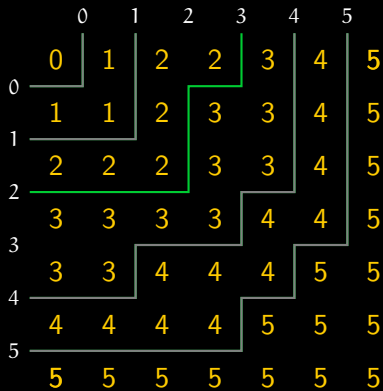
- $s_1, \dots, s_{n-r} \in V$ be a choice of “starting vertices”
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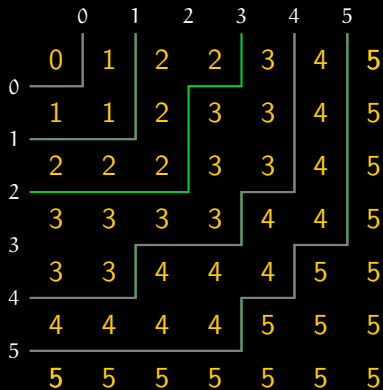
Then:

$$\boxed{\text{number of non-crossing path tuples from } (s_1, \dots, s_{n-r}) \text{ to } (e_1, \dots, e_{n-r})} = \det((a_{i,j}))_{i,j=1}^{n-r}.$$

0	1	2	2	3	4	5
1	1	2	3	3	4	5
2	2	2	3	3	4	5
3	3	3	3	4	4	5
3	3	4	4	4	5	5
4	4	4	4	5	5	5
5	5	5	5	5	5	5

0	1	2	2	3	4	5
1	1	2	3	3	4	5
2	2	2	3	3	4	5
3	3	3	3	4	4	5
3	3	4	4	4	5	5
4	4	4	4	5	5	5
5	5	5	5	5	5	5






$$\boxed{\# \text{paths from } u \text{ to } v} = \binom{u+v}{v}$$

$$\begin{vmatrix} \binom{0}{0} & \binom{1}{0} & \binom{2}{0} & \binom{3}{0} & \binom{4}{0} & \binom{5}{0} \\ \binom{1}{1} & \binom{2}{1} & \binom{3}{1} & \binom{4}{1} & \binom{5}{1} & \binom{6}{1} \\ \binom{2}{2} & \binom{3}{2} & \binom{4}{2} & \binom{5}{2} & \binom{6}{2} & \binom{7}{2} \\ \binom{3}{3} & \binom{4}{3} & \binom{5}{3} & \binom{6}{3} & \binom{7}{3} & \binom{8}{3} \\ \binom{4}{4} & \binom{5}{4} & \binom{6}{4} & \binom{7}{4} & \binom{8}{4} & \binom{9}{4} \\ \binom{5}{5} & \binom{6}{5} & \binom{7}{5} & \binom{8}{5} & \binom{9}{5} & \binom{10}{5} \end{vmatrix}$$

2 is not an
ending position



$\binom{0}{0}$	$\binom{1}{0}$	$\binom{2}{0}$	$\binom{3}{0}$	$\binom{4}{0}$	$\binom{5}{0}$
$\binom{1}{1}$	$\binom{2}{1}$	$\binom{3}{1}$	$\binom{4}{1}$	$\binom{5}{1}$	$\binom{6}{1}$
$\binom{2}{2}$	$\binom{3}{2}$	$\binom{4}{2}$	$\binom{5}{2}$	$\binom{6}{2}$	$\binom{7}{2}$
$\binom{3}{3}$	$\binom{4}{3}$	$\binom{5}{3}$	$\binom{6}{3}$	$\binom{7}{3}$	$\binom{8}{3}$
$\binom{4}{4}$	$\binom{5}{4}$	$\binom{6}{4}$	$\binom{7}{4}$	$\binom{8}{4}$	$\binom{9}{4}$
$\binom{5}{5}$	$\binom{6}{5}$	$\binom{7}{5}$	$\binom{8}{5}$	$\binom{9}{5}$	$\binom{10}{5}$

2 is not an
ending position

$\binom{0}{0}$	$\binom{1}{0}$	$\binom{2}{0}$	$\binom{3}{0}$	$\binom{4}{0}$	$\binom{5}{0}$
$\binom{1}{1}$	$\binom{2}{1}$	$\binom{3}{1}$	$\binom{4}{1}$	$\binom{5}{1}$	$\binom{6}{1}$
$\binom{2}{2}$	$\binom{3}{2}$	$\binom{4}{2}$	$\binom{5}{2}$	$\binom{6}{2}$	$\binom{7}{2}$
$\binom{3}{3}$	$\binom{4}{3}$	$\binom{5}{3}$	$\binom{6}{3}$	$\binom{7}{3}$	$\binom{8}{3}$
$\binom{4}{4}$	$\binom{5}{4}$	$\binom{6}{4}$	$\binom{7}{4}$	$\binom{8}{4}$	$\binom{9}{4}$
$\binom{5}{5}$	$\binom{6}{5}$	$\binom{7}{5}$	$\binom{8}{5}$	$\binom{9}{5}$	$\binom{10}{5}$

3 is not a
starting position

2 is not an
ending position

$$\Delta(5)_3^2 =$$

$\binom{0}{0}$	$\binom{1}{0}$	$\binom{2}{0}$	$\binom{3}{0}$	$\binom{4}{0}$	$\binom{5}{0}$
$\binom{1}{1}$	$\binom{2}{1}$	$\binom{3}{1}$	$\binom{4}{1}$	$\binom{5}{1}$	$\binom{6}{1}$
$\binom{2}{2}$	$\binom{3}{2}$	$\binom{4}{2}$	$\binom{5}{2}$	$\binom{6}{2}$	$\binom{7}{2}$
$\binom{3}{3}$	$\binom{4}{3}$	$\binom{5}{3}$	$\binom{6}{3}$	$\binom{7}{3}$	$\binom{8}{3}$
$\binom{4}{4}$	$\binom{5}{4}$	$\binom{6}{4}$	$\binom{7}{4}$	$\binom{8}{4}$	$\binom{9}{4}$
$\binom{5}{5}$	$\binom{6}{5}$	$\binom{7}{5}$	$\binom{8}{5}$	$\binom{9}{5}$	$\binom{10}{5}$

3 is not a
starting position

2 is not an
ending position

$$\Delta(5)_3^2 =$$

$\binom{0}{0}$	$\binom{1}{0}$	$\binom{2}{0}$	$\binom{3}{0}$	$\binom{4}{0}$	$\binom{5}{0}$
$\binom{1}{1}$	$\binom{2}{1}$	$\binom{3}{1}$	$\binom{4}{1}$	$\binom{5}{1}$	$\binom{6}{1}$
$\binom{2}{2}$	$\binom{3}{2}$	$\binom{4}{2}$	$\binom{5}{2}$	$\binom{6}{2}$	$\binom{7}{2}$
$\binom{3}{3}$	$\binom{4}{3}$	$\binom{5}{3}$	$\binom{6}{3}$	$\binom{7}{3}$	$\binom{8}{3}$
$\binom{4}{4}$	$\binom{5}{4}$	$\binom{6}{4}$	$\binom{7}{4}$	$\binom{8}{4}$	$\binom{9}{4}$
$\binom{5}{5}$	$\binom{6}{5}$	$\binom{7}{5}$	$\binom{8}{5}$	$\binom{9}{5}$	$\binom{10}{5}$

3 is not a
starting position

$$H_1(n, n) = \sum_{i=0}^{n-2} \sum_{j=0}^{n-2} \Delta(n-1)_i^j.$$

Homework: Show that

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- $\det\left(\binom{u+v}{v}\right)_{u,v=0}^{n-1} = 1$

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Homework: Show that

- $\det\left(\left(\binom{u+v}{v}\right)\right)_{u,v=0}^{n-1} = 1$
- $\Delta(n)_i^j = \sum_{\ell=0}^{n-1} \binom{\ell}{i} \binom{\ell}{j}$
- $\sum_{i=0}^{n-2} \sum_{j=0}^{n-2} \Delta(n-1)_i^j = \frac{1}{3}(4^{n-1} - 1)$

0	1	2	2	3	4	4
1	1	2	2	3	4	4
2	2	2	3	3	4	4
3	3	3	3	4	4	4
3	3	4	4	4	4	4
3	3	4	4	4	4	4
4	4	4	4	4	4	4

	0	1	2	3	4	5	
0	0	1	2	2	3	4	4
1	1	1	2	2	3	4	4
2	2	2	2	3	3	4	4
3	3	3	3	3	4	4	4
4	3	3	4	4	4	4	4
5	3	3	4	4	4	4	4
	4	4	4	4	4	4	4

2 is not an ending position 5 is not an ending position

$\binom{0}{0}$	$\binom{1}{0}$	$\binom{2}{0}$	$\binom{3}{0}$	$\binom{4}{0}$	$\binom{5}{0}$	
$\binom{1}{1}$	$\binom{2}{1}$	$\binom{3}{1}$	$\binom{4}{1}$	$\binom{5}{1}$	$\binom{6}{1}$	
$\binom{2}{2}$	$\binom{3}{2}$	$\binom{4}{2}$	$\binom{5}{2}$	$\binom{6}{2}$	$\binom{7}{2}$	
$\binom{3}{3}$	$\binom{4}{3}$	$\binom{5}{3}$	$\binom{6}{3}$	$\binom{7}{3}$	$\binom{8}{3}$	3 is not a starting position
$\binom{4}{4}$	$\binom{5}{4}$	$\binom{6}{4}$	$\binom{7}{4}$	$\binom{8}{4}$	$\binom{9}{4}$	4 is not a starting position
$\binom{5}{5}$	$\binom{6}{5}$	$\binom{7}{5}$	$\binom{8}{5}$	$\binom{9}{5}$	$\binom{10}{5}$	

$$\Delta(5)_{3,4}^{2,5} =$$

			2 is not an ending position		5 is not an ending position	
$\begin{pmatrix} 0 \\ 0 \end{pmatrix}$	$\begin{pmatrix} 1 \\ 0 \end{pmatrix}$	$\begin{pmatrix} 2 \\ 0 \end{pmatrix}$	$\begin{pmatrix} 3 \\ 0 \end{pmatrix}$	$\begin{pmatrix} 4 \\ 0 \end{pmatrix}$	$\begin{pmatrix} 5 \\ 0 \end{pmatrix}$	
$\begin{pmatrix} 1 \\ 1 \end{pmatrix}$	$\begin{pmatrix} 2 \\ 1 \end{pmatrix}$	$\begin{pmatrix} 3 \\ 1 \end{pmatrix}$	$\begin{pmatrix} 4 \\ 1 \end{pmatrix}$	$\begin{pmatrix} 5 \\ 1 \end{pmatrix}$	$\begin{pmatrix} 6 \\ 1 \end{pmatrix}$	
$\begin{pmatrix} 2 \\ 2 \end{pmatrix}$	$\begin{pmatrix} 3 \\ 2 \end{pmatrix}$	$\begin{pmatrix} 4 \\ 2 \end{pmatrix}$	$\begin{pmatrix} 5 \\ 2 \end{pmatrix}$	$\begin{pmatrix} 6 \\ 2 \end{pmatrix}$	$\begin{pmatrix} 7 \\ 2 \end{pmatrix}$	
$\begin{pmatrix} 3 \\ 3 \end{pmatrix}$	$\begin{pmatrix} 4 \\ 3 \end{pmatrix}$	$\begin{pmatrix} 5 \\ 3 \end{pmatrix}$	$\begin{pmatrix} 6 \\ 3 \end{pmatrix}$	$\begin{pmatrix} 7 \\ 3 \end{pmatrix}$	$\begin{pmatrix} 8 \\ 3 \end{pmatrix}$	3 is not a starting position
$\begin{pmatrix} 4 \\ 4 \end{pmatrix}$	$\begin{pmatrix} 5 \\ 4 \end{pmatrix}$	$\begin{pmatrix} 6 \\ 4 \end{pmatrix}$	$\begin{pmatrix} 7 \\ 4 \end{pmatrix}$	$\begin{pmatrix} 8 \\ 4 \end{pmatrix}$	$\begin{pmatrix} 9 \\ 4 \end{pmatrix}$	4 is not a starting position
$\begin{pmatrix} 5 \\ 5 \end{pmatrix}$	$\begin{pmatrix} 6 \\ 5 \end{pmatrix}$	$\begin{pmatrix} 7 \\ 5 \end{pmatrix}$	$\begin{pmatrix} 8 \\ 5 \end{pmatrix}$	$\begin{pmatrix} 9 \\ 5 \end{pmatrix}$	$\begin{pmatrix} 10 \\ 5 \end{pmatrix}$	

$$\Delta(5)_{3,4}^{2,5} =$$

			2 is not an ending position		5 is not an ending position	
$\binom{0}{0}$	$\binom{1}{0}$	$\binom{2}{0}$	$\binom{3}{0}$	$\binom{4}{0}$	$\binom{5}{0}$	
$\binom{1}{1}$	$\binom{2}{1}$	$\binom{3}{1}$	$\binom{4}{1}$	$\binom{5}{1}$	$\binom{6}{1}$	
$\binom{2}{2}$	$\binom{3}{2}$	$\binom{4}{2}$	$\binom{5}{2}$	$\binom{6}{2}$	$\binom{7}{2}$	
$\binom{3}{3}$	$\binom{4}{3}$	$\binom{5}{3}$	$\binom{6}{3}$	$\binom{7}{3}$	$\binom{8}{3}$	3 is not a starting position
$\binom{4}{4}$	$\binom{5}{4}$	$\binom{6}{4}$	$\binom{7}{4}$	$\binom{8}{4}$	$\binom{9}{4}$	4 is not a starting position
$\binom{5}{5}$	$\binom{6}{5}$	$\binom{7}{5}$	$\binom{8}{5}$	$\binom{9}{5}$	$\binom{10}{5}$	

$$H_2(n, n) = \sum_{i_1=0}^{n-2} \sum_{i_2=i_1+1}^{n-2} \sum_{j_1=0}^{n-2} \sum_{j_2=j_1+1}^{n-2} \Delta(n-1)_{i_1, i_2}^{j_1, j_2}.$$

$$H_r(n, n) = \sum_{\substack{0 \leq i_1 < \dots < i_r < n-1 \\ 0 \leq j_1 < \dots < j_r < n-1}} \Delta(n-1)_{i_1, \dots, i_r}^{j_1, \dots, j_r}$$

$$H_r(n, n) = \sum_{\substack{0 \leq i_1 < \dots < i_r < n-1 \\ 0 \leq j_1 < \dots < j_r < n-1}} \Delta(n-1)_{i_1, \dots, i_r}^{j_1, \dots, j_r}$$

Is this D-finite?

Jacobi's determinant identity implies

$$\Delta(\mathbf{n}-1)_{i_1, \dots, i_r}^{j_1, \dots, j_r} = \begin{vmatrix} \Delta(\mathbf{n}-1)_{i_1}^{j_1} & \cdots & \Delta(\mathbf{n}-1)_{i_1}^{j_r} \\ \vdots & \ddots & \vdots \\ \Delta(\mathbf{n}-1)_{i_r}^{j_1} & \cdots & \Delta(\mathbf{n}-1)_{i_r}^{j_r} \end{vmatrix}$$

Jacobi's determinant identity implies

$$\Delta(n-1)_{i_1, \dots, i_r}^{j_1, \dots, j_r} = \begin{vmatrix} \Delta(n-1)_{i_1}^{j_1} & \cdots & \Delta(n-1)_{i_1}^{j_r} \\ \vdots & \ddots & \vdots \\ \Delta(n-1)_{i_r}^{j_1} & \cdots & \Delta(n-1)_{i_r}^{j_r} \end{vmatrix}$$

This is D-finite in n for every fixed r and $i_1, \dots, i_r, j_1, \dots, j_r$.

Jacobi's determinant identity implies

$$\Delta(n-1)_{i_1, \dots, i_r}^{j_1, \dots, j_r} = \begin{vmatrix} \Delta(n-1)_{i_1}^{j_1} & \cdots & \Delta(n-1)_{i_1}^{j_r} \\ \vdots & \ddots & \vdots \\ \Delta(n-1)_{i_r}^{j_1} & \cdots & \Delta(n-1)_{i_r}^{j_r} \end{vmatrix}$$

This is D-finite in n for every fixed r and $i_1, \dots, i_r, j_1, \dots, j_r$.

Therefore, $H_r(n, n) = \sum_{i_1, \dots, i_r} \sum_{j_1, \dots, j_r} \Delta(n-1)_{i_1, \dots, i_r}^{j_1, \dots, j_r}$ is D-finite.

$$\begin{aligned}
& (201600n^9 + 4942080n^8 + 53078112n^7 + 327661728n^6 + \\
& 1280700480n^5 + 3285342016n^4 + 5528828352n^3 + 5883447104n^2 + \\
& 3591093120n + 957662208)a(n) + (-970200n^9 - 24199560n^8 - \\
& 264810744n^7 - 1667830872n^6 - 6659340648n^5 - 17470825688n^4 - \\
& 30096410912n^3 - 32804461872n^2 - 20514211488n - \\
& 5603970816)a(n+1) + (589050n^9 + 14827590n^8 + 163756656n^7 + \\
& 1040895564n^6 + 4194035058n^5 + 11101344742n^4 + 19289250308n^3 + \\
& 21198776056n^2 + 13360158000n + 3676219776)a(n+2) + (294525n^9 + \\
& 7319295n^8 + 79828578n^7 + 501335472n^6 + 1997003589n^5 + \\
& 5229549731n^4 + 8997110634n^3 + 9799013608n^2 + 6125859120n + \\
& 1673566848)a(n+3) + (-121275n^9 - 3053295n^8 - 33716268n^7 - \\
& 214212552n^6 - 862421763n^5 - 2280190003n^4 - 3956305720n^3 - \\
& 4340670060n^2 - 2730542400n - 749859264)a(n+4) + (6300n^9 + \\
& 163890n^8 + 1863666n^7 + 12150660n^6 + 50023284n^5 + 134779202n^4 + \\
& 237527338n^3 + 263895164n^2 + 167643648n + 46381248)a(n+5) \stackrel{?}{=} 0
\end{aligned}$$

$$H_2(n, n) = \sum_{i_1=0}^{n-2} \sum_{i_2=i_1+1}^{n-2} \sum_{j_1=0}^{n-2} \sum_{j_2=j_1+1}^{n-2} \Delta(n-1)_{i_1, i_2}^{j_1, j_2}$$

$$H_2(n, n) = \sum_{i_1=0}^{n-2} \sum_{i_2=i_1+1}^{n-2} \sum_{j_1=0}^{n-2} \sum_{j_2=j_1+1}^{n-2} \begin{vmatrix} \Delta(n-1)_{i_1}^{j_1} & \Delta(n-1)_{i_1}^{j_2} \\ \Delta(n-1)_{i_2}^{j_1} & \Delta(n-1)_{i_2}^{j_2} \end{vmatrix}$$

$$H_2(n, n) = \sum_{i_1=0}^{n-2} \sum_{i_2=i_1+1}^{n-2} \sum_{j_1=0}^{n-2} \sum_{j_2=j_1+1}^{n-2} \left| \begin{array}{cc} \sum_{\ell=0}^{n-1} \binom{\ell}{i_1} \binom{\ell}{j_1} & \sum_{\ell=0}^{n-1} \binom{\ell}{i_1} \binom{\ell}{j_2} \\ \sum_{\ell=0}^{n-1} \binom{\ell}{i_2} \binom{\ell}{j_1} & \sum_{\ell=0}^{n-1} \binom{\ell}{i_2} \binom{\ell}{j_2} \end{array} \right|$$

$$H_2(n, n) = S_1(n) - S_2(n), \text{ where}$$

$H_2(n, n) = S_1(n) - S_2(n)$, where

$$S_1(n) = \sum_{i_1 \geq 0} \sum_{i_2 > i_1} \sum_{j_1 \geq 0} \sum_{j_2 > j_1} \sum_{u=0}^n \sum_{v=0}^n \binom{u}{i_1} \binom{u}{j_1} \binom{v}{i_2} \binom{v}{j_2}$$

$$S_2(n) = \sum_{i_1 \geq 0} \sum_{i_2 > i_1} \sum_{j_1 \geq 0} \sum_{j_2 > j_1} \sum_{u=0}^n \sum_{v=0}^n \binom{u}{i_1} \binom{u}{j_2} \binom{v}{i_2} \binom{v}{j_1}$$

$H_2(n, n) = S_1(n) - S_2(n)$, where

$$S_1(n) = \sum_{u=0}^n \sum_{v=0}^n \left(\sum_{i_1 \geq 0} \sum_{i_2 > i_1} \binom{u}{i_1} \binom{v}{i_2} \right) \left(\sum_{j_1 \geq 0} \sum_{j_2 > j_1} \binom{u}{j_1} \binom{v}{j_2} \right)$$

$$S_2(n) = \sum_{u=0}^n \sum_{v=0}^n \left(\sum_{i_1 \geq 0} \sum_{i_2 > i_1} \binom{u}{i_1} \binom{v}{i_2} \right) \left(\sum_{j_1 \geq 0} \sum_{j_2 > j_1} \binom{v}{j_1} \binom{u}{j_2} \right)$$

$H_2(n, n) = S_1(n) - S_2(n)$, where

$$S_1(n) = \sum_{u=0}^n \sum_{v=0}^n \underbrace{\left(\sum_{i_1 \geq 0} \sum_{i_2 > i_1} \binom{u}{i_1} \binom{v}{i_2} \right)}_{=:s(u,v)} \underbrace{\left(\sum_{j_1 \geq 0} \sum_{j_2 > j_1} \binom{u}{j_1} \binom{v}{j_2} \right)}_{=:s(u,v)}$$

$$S_2(n) = \sum_{u=0}^n \sum_{v=0}^n \underbrace{\left(\sum_{i_1 \geq 0} \sum_{i_2 > i_1} \binom{u}{i_1} \binom{v}{i_2} \right)}_{=:s(u,v)} \underbrace{\left(\sum_{j_1 \geq 0} \sum_{j_2 > j_1} \binom{v}{j_1} \binom{u}{j_2} \right)}_{=:s(v,u)}$$

$H_2(n, n) = S_1(n) - S_2(n)$, where

$$S_1(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)^2$$

$$S_2(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)s(v, u)$$

$$S_1(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)^2$$

$$S_2(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)s(v, u)$$

$$S_1(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)^2 \qquad S_2(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)s(v, u)$$

$$\sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u, v) x^u y^v$$

$$S_1(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)^2 \quad S_2(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)s(v, u)$$

$$\sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u, v) x^u y^v = \frac{y}{(1-x-y)(1-2y)}$$

$$S_1(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)^2 \quad S_2(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)s(v, u)$$

$$f(x, y) := \sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u, v) x^u y^v = \frac{y}{(1-x-y)(1-2y)}$$

$$S_1(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)^2 \quad S_2(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v) s(v, u)$$

$$f(x, y) := \sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u, v) x^u y^v = \frac{y}{(1-x-y)(1-2y)}$$

$$\sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u, v)^2 x^u y^v = f(x, y) \odot_{x, y} f(x, y)$$

$$\sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u, v) s(v, u) x^u y^v = f(x, y) \odot_{x, y} f(y, x)$$

$$S_1(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)^2 \quad S_2(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v) s(v, u)$$

$$f(x, y) := \sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u, v) x^u y^v = \frac{y}{(1-x-y)(1-2y)}$$

$$\sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u, v)^2 x^u y^v = [X^{-1}][Y^{-1}] \frac{1}{XY} f\left(\frac{x}{X}, \frac{y}{Y}\right) f(X, Y)$$

$$\sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u, v) s(v, u) x^u y^v = [X^{-1}][Y^{-1}] \frac{1}{XY} f\left(\frac{x}{X}, \frac{y}{Y}\right) f(Y, X)$$

$$S_1(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)^2 \quad S_2(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)s(v, u)$$

$$f(x, y) := \sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u, v) x^u y^v = \frac{y}{(1-x-y)(1-2y)}$$

$$\sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u, v)^2 x^u y^v = g_1(x, y)$$

$$\sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u, v)s(v, u) x^u y^v = g_2(x, y)$$

$$S_1(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)^2 \quad S_2(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)s(v, u)$$

$$f(x, y) := \sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u, v) x^u y^v = \frac{y}{(1-x-y)(1-2y)}$$

$$\frac{1}{(1-x)(1-y)} \sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u, v)^2 x^u y^v = \frac{1}{(1-x)(1-y)} g_1(x, y)$$

$$\frac{1}{(1-x)(1-y)} \sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u, v)s(v, u) x^u y^v = \frac{1}{(1-x)(1-y)} g_2(x, y)$$

$$S_1(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)^2 \quad S_2(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v) s(v, u)$$

$$f(x, y) := \sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u, v) x^u y^v = \frac{y}{(1-x-y)(1-2y)}$$

$$\sum_{n=0}^{\infty} \sum_{m=0}^{\infty} \sum_{u=0}^n \sum_{v=0}^m s(u, v)^2 x^n y^m = \frac{1}{(1-x)(1-y)} g_1(x, y)$$

$$\sum_{n=0}^{\infty} \sum_{m=0}^{\infty} \sum_{u=0}^n \sum_{v=0}^m s(u, v) s(v, u) x^n y^m = \frac{1}{(1-x)(1-y)} g_2(x, y)$$

$$S_1(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)^2 \quad S_2(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)s(v, u)$$

$$f(x, y) := \sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u, v) x^u y^v = \frac{y}{(1-x-y)(1-2y)}$$

$$\sum_{n=0}^{\infty} \sum_{m=0}^{\infty} \sum_{u=0}^n \sum_{v=0}^m s(u, v)^2 x^n y^m = \frac{1}{(1-x)(1-y)} g_1(x, y)$$

$$\sum_{n=0}^{\infty} \sum_{m=0}^{\infty} \sum_{u=0}^n \sum_{v=0}^m s(u, v)s(v, u) x^n y^m = \frac{1}{(1-x)(1-y)} g_2(x, y)$$

$$S_1(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)^2 \quad S_2(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v) s(v, u)$$

$$f(x, y) := \sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u, v) x^u y^v = \frac{y}{(1-x-y)(1-2y)}$$

$$\sum_{n=0}^{\infty} \sum_{u=0}^n \sum_{v=0}^n s(u, v)^2 x^n = [y^{-1}] \frac{1}{y(1-\frac{x}{y})(1-y)} g_1(\frac{x}{y}, y)$$

$$\sum_{n=0}^{\infty} \sum_{u=0}^n \sum_{v=0}^n s(u, v) s(v, u) x^n = [y^{-1}] \frac{1}{y(1-\frac{x}{y})(1-y)} g_2(\frac{x}{y}, y)$$

$$S_1(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v)^2 \quad S_2(n) = \sum_{u=0}^n \sum_{v=0}^n s(u, v) s(v, u)$$

$$f(x, y) := \sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u, v) x^u y^v = \frac{y}{(1-x-y)(1-2y)}$$

$$\sum_{n=0}^{\infty} S_1(n) x^n = [y^{-1}] \frac{1}{y(1-\frac{x}{y})(1-y)} g_1(\frac{x}{y}, y)$$

$$\sum_{n=0}^{\infty} S_2(n) x^n = [y^{-1}] \frac{1}{y(1-\frac{x}{y})(1-y)} g_2(\frac{x}{y}, y)$$

Altogether,

$$\sum_{n=0}^{\infty} H_2(n) x^n = [y^{-1}] \frac{1}{y(1-\frac{x}{y})(1-y)} \left([X^{-1}][Y^{-1}] \frac{1}{XY} f\left(\frac{x}{Xy}, \frac{y}{Y}\right) f(X, Y) \right. \\ \left. - [X^{-1}][Y^{-1}] \frac{1}{XY} f\left(\frac{x}{Xy}, \frac{y}{Y}\right) f(Y, X) \right)$$

Altogether,

$$\sum_{n=0}^{\infty} H_2(n) x^n = \underset{\substack{\uparrow \\ \text{diagonal}}}{[y^{-1}]} \frac{1}{y(1-\frac{x}{y})(1-y)} \left(\underset{\substack{\uparrow \\ \text{Hadamard} \\ \text{product}}}{[X^{-1}][Y^{-1}]} \frac{1}{XY} f\left(\frac{x}{Xy}, \frac{y}{Y}\right) f(X, Y) \right. \\ \left. - [X^{-1}][Y^{-1}] \frac{1}{XY} f\left(\frac{x}{Xy}, \frac{y}{Y}\right) f(Y, X) \right)$$

Altogether,

$$\sum_{n=0}^{\infty} H_2(n) x^n = \underset{\substack{\uparrow \\ \text{diagonal}}}{[y^{-1}]} \frac{1}{y(1-\frac{x}{y})(1-y)} \left(\underset{\substack{\uparrow \\ \text{Hadamard} \\ \text{product}}}{[X^{-1}][Y^{-1}]} \frac{1}{XY} f\left(\frac{x}{Xy}, \frac{y}{Y}\right) f(X, Y) \right. \\ \left. - [X^{-1}][Y^{-1}] \frac{1}{XY} f\left(\frac{x}{Xy}, \frac{y}{Y}\right) f(Y, X) \right)$$

From this expression we can compute a recurrence for $H_2(n)$.

Altogether,

$$\sum_{n=0}^{\infty} H_2(n)x^n = \underset{\substack{\uparrow \\ \text{diagonal}}}{[y^{-1}]} \frac{1}{y(1-\frac{x}{y})(1-y)} \left(\underset{\substack{\uparrow \\ \text{Hadamard} \\ \text{product}}}{[X^{-1}][Y^{-1}]} \frac{1}{XY} f\left(\frac{x}{Xy}, \frac{y}{Y}\right) f(X, Y) \right. \\ \left. - \underset{\substack{\uparrow \\ \text{Hadamard} \\ \text{product}}}{[X^{-1}][Y^{-1}]} \frac{1}{XY} f\left(\frac{x}{Xy}, \frac{y}{Y}\right) f(Y, X) \right)$$

From this expression we can compute a recurrence for $H_2(n)$.

This proves the conjecture.

$$H_3(n, n) = \sum_{i_1=0}^{n-2} \sum_{i_2=i_1+1}^{n-2} \sum_{i_3=i_2+1}^{n-2} \sum_{j_1=0}^{n-2} \sum_{j_2=j_1+1}^{n-2} \sum_{j_3=j_2+1}^{n+2} \Delta(n-1)_{i_1, i_2, i_3}^{j_1, j_2, j_3}$$

$$H_3(n, n) = \sum_{i_1=0}^{n-2} \sum_{i_2=i_1+1}^{n-2} \sum_{i_3=i_2+1}^{n-2} \sum_{j_1=0}^{n-2} \sum_{j_2=j_1+1}^{n-2} \sum_{j_3=j_2+1}^{n+2} \Delta(n-1)_{i_1, i_2, i_3}^{j_1, j_2, j_3}$$

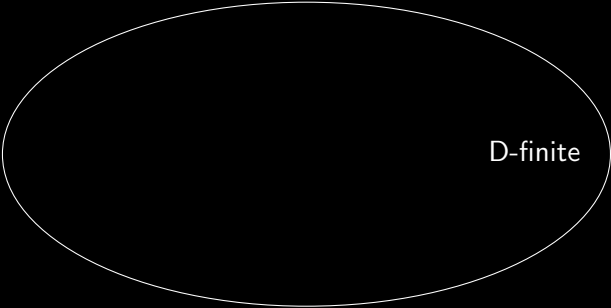
For this case, we only have a guessed recurrence.



What's next?

Papers related to D-finiteness at this year's ISSAC:

- **Louis Gaillard:** A unified approach for degree bound estimates of linear differential operators
- **Shaoshi Chen, Manuel Kauers, Christoph Koutschan, Xiuyun Li, Rong-Hua Wang and Yisen Wang:** Non-minimality of minimal telescopers explained by residues
- **Manuel Kauers and Raphael Pages:** Bounds for D-Algebraic Closure Properties
- **Alaa Ibrahim:** Positivity Proofs for Linear Recurrences with Several Dominant Eigenvalues
- **Jérémy Berthomieu, Romain Lebreton and Kevin Tran:** Quasi-Linear Guessing of Minimal Lexicographic Gröbner Bases of Ideals of C-Relations of Random Bi-Indexed Sequences



D-finite

