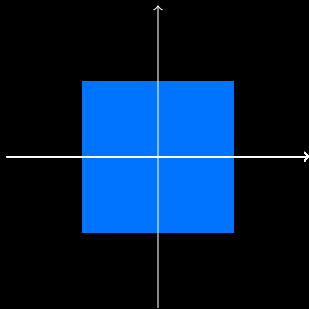


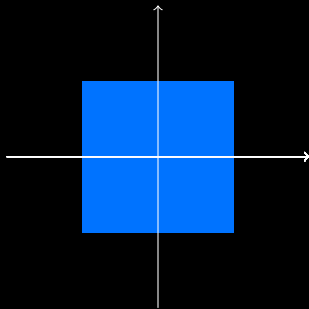
SYMMETRY BREAKING FOR QUANTIFIED BOOLEAN FORMULAS



Manuel Kauers · JKU

Joint work with Martina Seidl (SAT'18)





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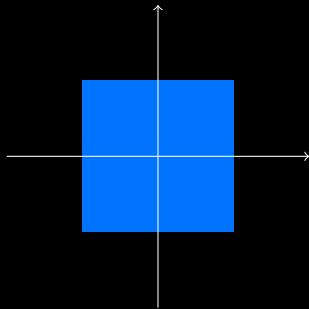
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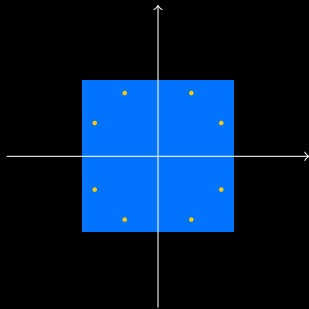
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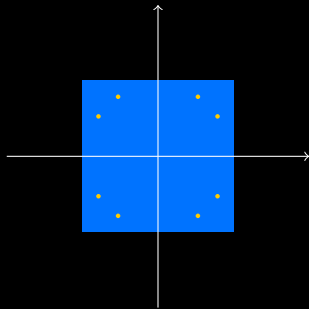
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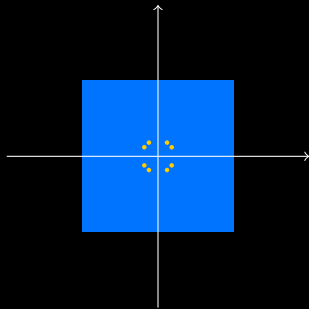
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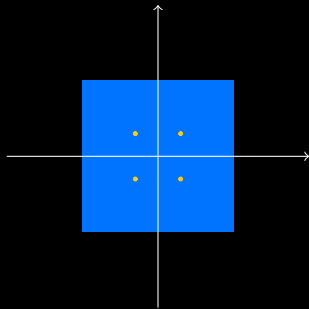
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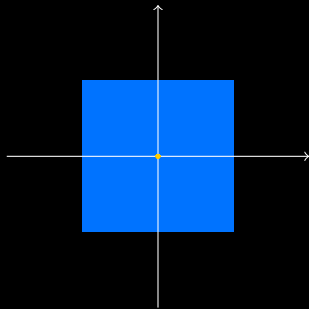


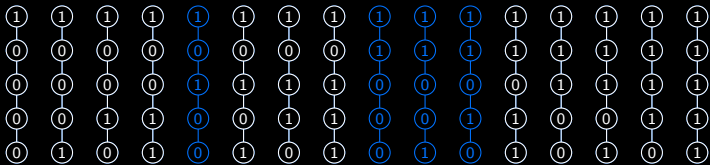
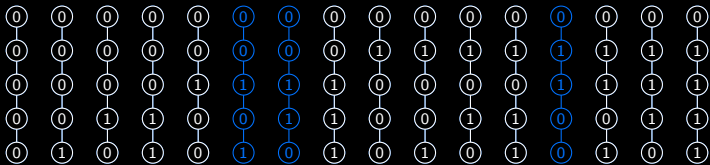












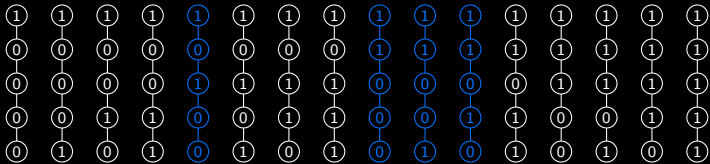
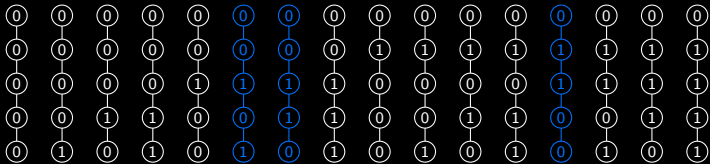
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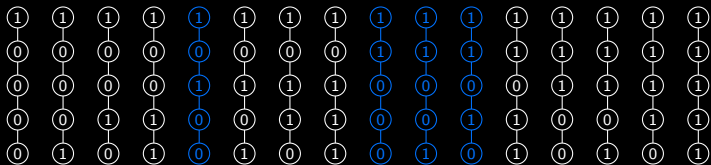
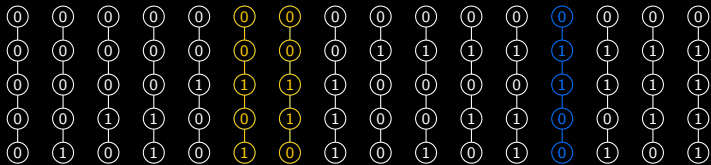
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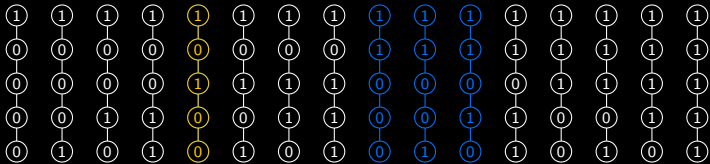
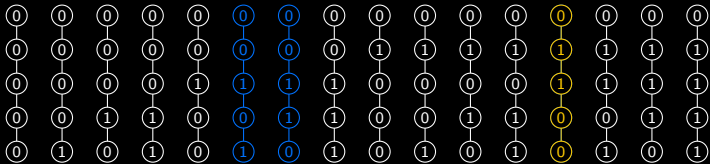
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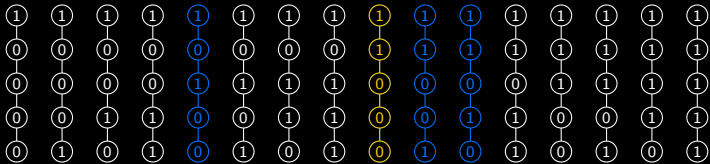
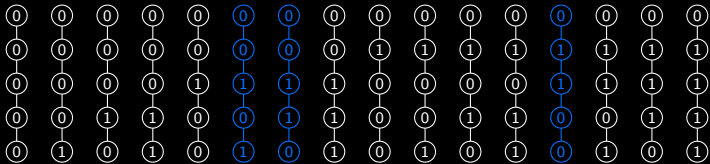
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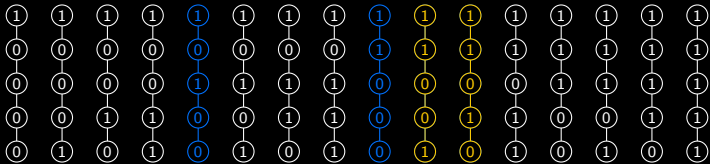
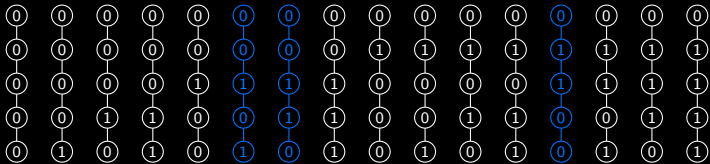
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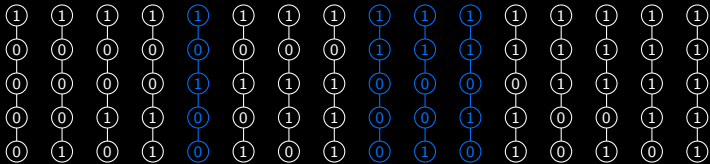
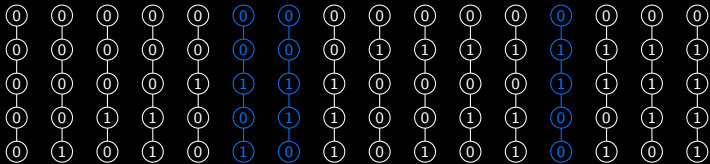


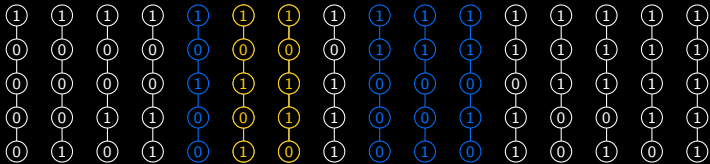
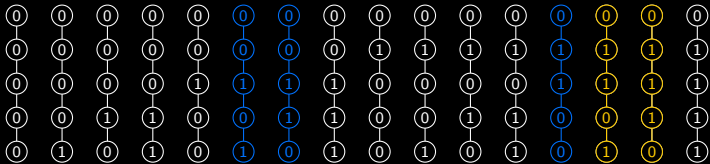












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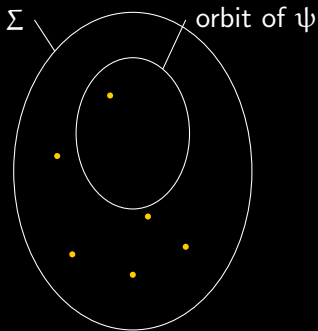
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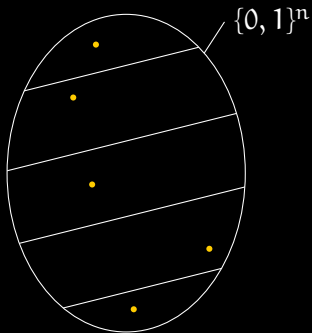
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syntactic



every $\sigma \in \{0, 1\}^n$ is
a solution of at least one
formula in the orbit of ψ

semantic



every orbit of $\{0, 1\}^n$
contains at least one
solution of ψ

For SAT, there is no difference if we restrict to “easy” symmetries.

$$\phi(x_1 , x_2 , x_3 , x_4 , x_5)$$



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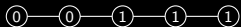
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- Observe that we use a “syntactic” group but a “semantic” justification.

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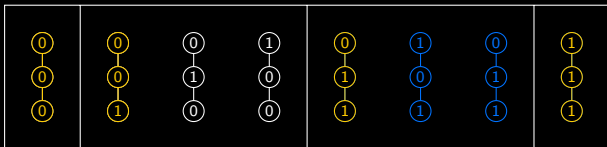
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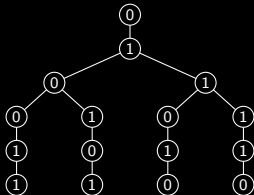
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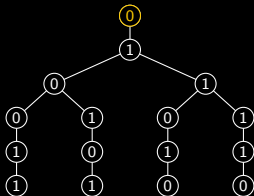
What about QBF?

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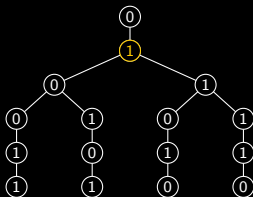
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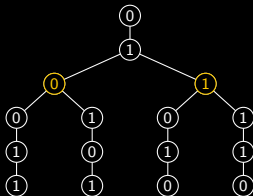
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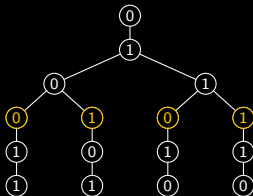
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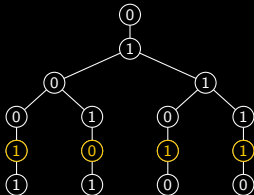
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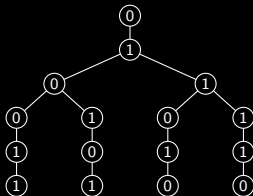
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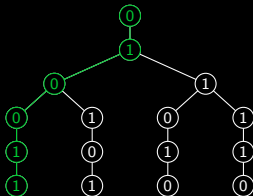
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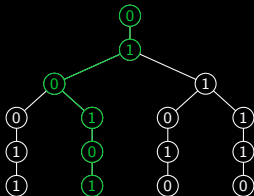
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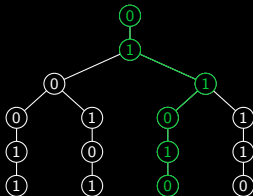
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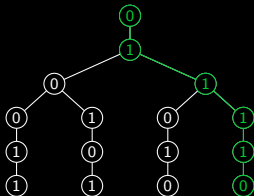
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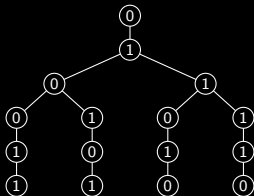
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Given a quantifier prefix P , we write $\mathbb{S}(P)$ for the corresponding set of tree assignments.

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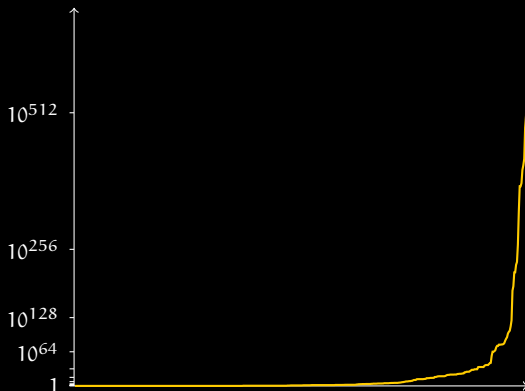
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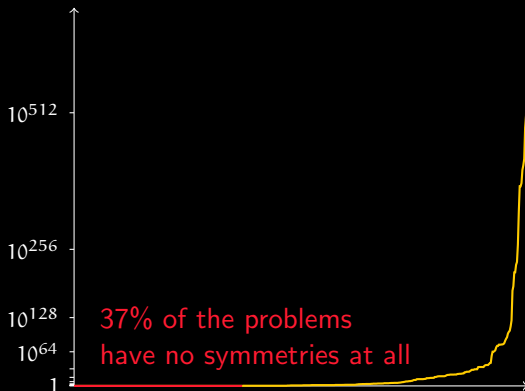
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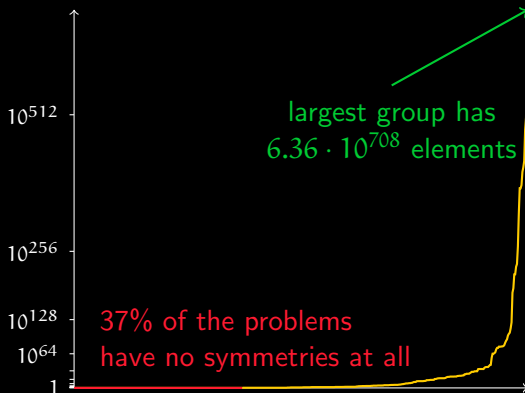
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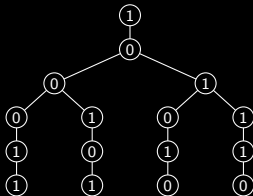
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- **Key fact:** If G_{syn} is a syntactic symmetry group for $P.\phi$ and G_{sem} is a semantic symmetry group for $P.\phi$ and ψ is a symmetry breaker for G_{syn} and G_{sem} , then $P.\phi$ has a solution in $\mathbb{S}(P)$ if and only if $P.(\phi \wedge \psi)$ does.

Easy symmetries

What does permutation of variables mean semantically?

$$\exists x_1, x_2 \forall x_3, x_4 \exists x_5, x_6 : \phi(x_1, x_2, x_3, x_4, x_5, x_6)$$



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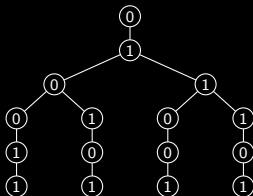
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Unlike in SAT, there is no longer a 1:1 correspondence.

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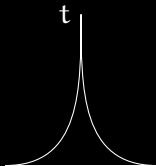
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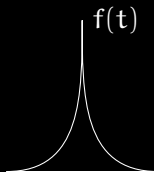
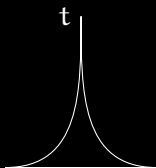
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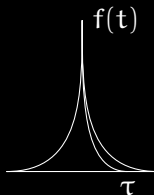
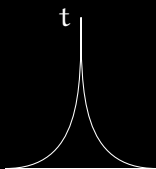
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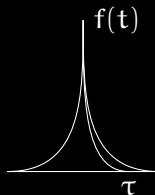
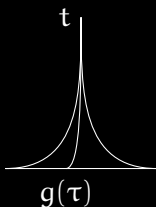
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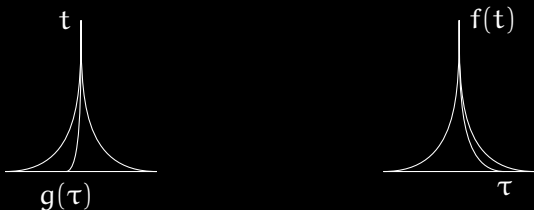
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- Then G_{sem} is called the **associated group** for G_{syn} .

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- Observe that only G_{syn} appears in the formula. The group G_{sem} is only used for the justification.

Case study: the Kleine-Büning et al family

Solving times of DepQBF (in sec)

w/o SB

with SB

n	QRes	LD	QRes	LD
10	0.3	0.5	0.4	0.4
20	160	0.5	0.4	0.4
40	> 3600	0.5	0.4	0.4
80	> 3600	0.7	0.4	0.4
160	> 3600	2.2	0.5	0.4
320	> 3600	12.3	0.6	0.5
640	> 3600	36.8	1.0	0.8
1280	> 3600	241.1	22.6	19.7
2560	> 3600	> 3600	215.7	155.2
5120	> 3600	> 3600	1873.2	1042.6

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In fact, we can show that there
is a proof of size $O(n)$
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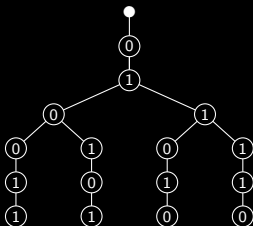
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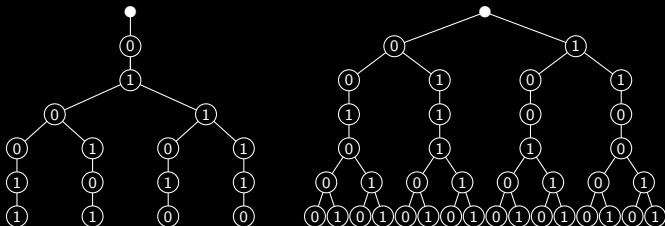
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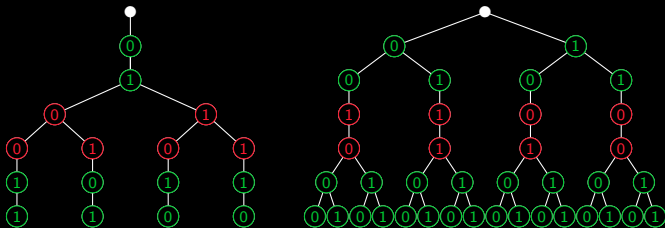
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$$P.\phi \text{ is true} \iff P.((\phi \wedge \psi_{\exists}) \vee \psi_{\forall}) \text{ is true}$$

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