

# THE FBHHRBNRSSHK-ALGORITHM FOR MULTIPLICATION IN $\mathbb{Z}_2^{5 \times 5}$ IS STILL NOT THE END OF THE STORY

MANUEL KAUERS\* AND JAKOB MOOSBAUER†

ABSTRACT. In response to a recent *Nature* article which announced an algorithm for multiplying  $5 \times 5$ -matrices over  $\mathbb{Z}_2$  with only 96 multiplications, two fewer than the previous record, we present an algorithm that does the job with only 95 multiplications.

## 1. INTRODUCTION

Ever since Strassen [8] discovered that  $2 \times 2$ -matrices can be multiplied with only 7 multiplications in the coefficient domain, there is a mystery around the complexity of matrix multiplication. For asymptotically large  $n$ , the best we know at the moment is a multiplication algorithm that requires  $O(n^{2.3728596})$  operations [1], slightly improving upon the previous record  $O(n^{2.3728639})$  [5]. For  $n = 3$ , it is known that 23 multiplications suffice in a non-commutative setting [4]. For  $n = 4$ , we can solve the problem with 49 multiplications by applying Strassen's algorithm recursively. In a recent article that received considerable media attention, Fawzi et al. [2] used a machine learning approach to find a multiplication scheme with 47 multiplications, applicable to coefficient domains of characteristic 2. Under the same restriction on the coefficient domain, they also improved the best known bound for  $n = 5$  from 98 [7] to 96. See [2, 6] for current records for other formats and further references on the matter.

In this short note, we present another non-equivalent solution for  $4 \times 4$ -matrices requiring 47 multiplications, as well as a first solution for  $5 \times 5$ -matrices requiring 95 multiplications. This solution was obtained from the scheme of Fawzi et al. by applying a sequence of transformations leading to a scheme from which one multiplication could be eliminated. Our new scheme for  $4 \times 4$ -matrices was obtained by the same technique, taking the standard multiplication algorithm (with 64 multiplications) as starting point. We will explain our search technique in further detail in a forthcoming paper [3].

## 2. MULTIPLICATION OF $4 \times 4$ MATRICES

Let

$$A = \begin{pmatrix} a_{1,1} & a_{1,2} & a_{1,3} & a_{1,4} \\ a_{2,1} & a_{2,2} & a_{2,3} & a_{2,4} \\ a_{3,1} & a_{3,2} & a_{3,3} & a_{3,4} \\ a_{4,1} & a_{4,2} & a_{4,3} & a_{4,4} \end{pmatrix} \text{ and } B = \begin{pmatrix} b_{1,1} & b_{1,2} & b_{1,3} & b_{1,4} \\ b_{2,1} & b_{2,2} & b_{2,3} & b_{2,4} \\ b_{3,1} & b_{3,2} & b_{3,3} & b_{3,4} \\ b_{4,1} & b_{4,2} & b_{4,3} & b_{4,4} \end{pmatrix},$$

and

$$C = \begin{pmatrix} c_{1,1} & c_{1,2} & c_{1,3} & c_{1,4} \\ c_{2,1} & c_{2,2} & c_{2,3} & c_{2,4} \\ c_{3,1} & c_{3,2} & c_{3,3} & c_{3,4} \\ c_{4,1} & c_{4,2} & c_{4,3} & c_{4,4} \end{pmatrix} := AB$$

be matrices with coefficients in a ring  $R$  of characteristic 2.

We can then compute the entries of  $C$  from the entries of  $A$  and  $B$  by first setting

$$m_k = \left( \sum_{i=1}^4 \sum_{j=1}^4 \alpha_{i,j}^{(k)} a_{i,j} \right) \left( \sum_{i=1}^4 \sum_{j=1}^4 \beta_{i,j}^{(k)} b_{i,j} \right)$$

for  $k = 1, \dots, 47$  and then

$$c_{i,j} = \sum_{k=1}^{47} \gamma_{i,j}^{(k)} m_k$$

for  $i, j = 1, \dots, 4$ , where the  $\alpha_{i,j}^{(k)}, \beta_{i,j}^{(k)}, \gamma_{i,j}^{(k)}$  are as follows:

\* Supported by the Austrian FWF grant P31571-N32.

† Supported by the Land Oberösterreich through the Lit-AI Lab.







43	$\begin{pmatrix} 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix}$
44	$\begin{pmatrix} 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$
45	$\begin{pmatrix} 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 \end{pmatrix}$	$\begin{pmatrix} 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 \end{pmatrix}$	$\begin{pmatrix} 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix}$
46	$\begin{pmatrix} 1 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 1 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix}$
47	$\begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$	$\begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$	$\begin{pmatrix} 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix}$

An electronic version of this multiplication scheme is available at <http://www.algebra.uni-linz.ac.at/people/mkauers/matrix-mult/s47.exp>

### 3. MULTIPLICATION OF $5 \times 5$ MATRICES

Let

$$A = \begin{pmatrix} a_{1,1} & a_{1,2} & a_{1,3} & a_{1,4} & a_{1,5} \\ a_{2,1} & a_{2,2} & a_{2,3} & a_{2,4} & a_{2,5} \\ a_{3,1} & a_{3,2} & a_{3,3} & a_{3,4} & a_{3,5} \\ a_{4,1} & a_{4,2} & a_{4,3} & a_{4,4} & a_{4,5} \\ a_{5,1} & a_{5,2} & a_{5,3} & a_{5,4} & a_{5,5} \end{pmatrix} \text{ and } B = \begin{pmatrix} b_{1,1} & b_{1,2} & b_{1,3} & b_{1,4} & b_{1,5} \\ b_{2,1} & b_{2,2} & b_{2,3} & b_{2,4} & b_{2,5} \\ b_{3,1} & b_{3,2} & b_{3,3} & b_{3,4} & b_{3,5} \\ b_{4,1} & b_{4,2} & b_{4,3} & b_{4,4} & b_{4,5} \\ b_{5,1} & b_{5,2} & b_{5,3} & b_{5,4} & b_{5,5} \end{pmatrix},$$

and

$$C = \begin{pmatrix} c_{1,1} & c_{1,2} & c_{1,3} & c_{1,4} & c_{1,5} \\ c_{2,1} & c_{2,2} & c_{2,3} & c_{2,4} & c_{2,5} \\ c_{3,1} & c_{3,2} & c_{3,3} & c_{3,4} & c_{3,5} \\ c_{4,1} & c_{4,2} & c_{4,3} & c_{4,4} & c_{4,5} \\ c_{5,1} & c_{5,2} & c_{5,3} & c_{5,4} & c_{5,5} \end{pmatrix} := AB$$

be matrices with coefficients in a ring  $R$  of characteristic 2.

We can then compute the entries of  $C$  from the entries of  $A$  and  $B$  by first setting

$$m_k = \left( \sum_{i=1}^5 \sum_{j=1}^5 \alpha_{i,j}^{(k)} a_{i,j} \right) \left( \sum_{i=1}^5 \sum_{j=1}^5 \beta_{i,j}^{(k)} b_{i,j} \right)$$

for  $k = 1, \dots, 95$  and then

$$c_{i,j} = \sum_{k=1}^{95} \gamma_{i,j}^{(k)} m_k$$

for  $i, j = 1, \dots, 5$ , where the  $\alpha_{i,j}^{(k)}, \beta_{i,j}^{(k)}, \gamma_{i,j}^{(k)}$  are as follows:

$k$	$((\alpha_{i,j}^{(k)}))_{i,j=1}^5$	$((\beta_{i,j}^{(k)}))_{i,j=1}^5$	$((\gamma_{i,j}^{(k)}))_{i,j=1}^5$
1	$\begin{pmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}$



















90	$\begin{pmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{pmatrix}$
91	$\begin{pmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \end{pmatrix}$	$\begin{pmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \end{pmatrix}$	$\begin{pmatrix} 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \end{pmatrix}$
92	$\begin{pmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 1 & 1 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 \end{pmatrix}$
93	$\begin{pmatrix} 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}$
94	$\begin{pmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{pmatrix}$
95	$\begin{pmatrix} 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 \end{pmatrix}$	$\begin{pmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}$

An electronic version of this multiplication scheme is available at <http://www.algebra.uni-linz.ac.at/people/mkauers/matrix-mult/s95.exp>

#### REFERENCES

- [1] Josh Alman and Virginia Vassilevska Williams. A refined laser method and faster matrix multiplication. In *Proceedings of SODA'21*, page ArXiv 2010.05846, 2021.
- [2] Alhussein Fawzi, Matej Balog, Aja Huang, Thomas Hubert, Bernardino Romera-Paredes, Mohammadamin Barekatin, Alexander Novikov, Francisco J. R. Ruiz, Julian Schrittwieser, Grzegorz Swirszcz, David Silver, Demis Hassabis, and Pushmeet Kohli. Discovering faster matrix multiplication algorithms with reinforcement learning. *Nature*, 610:47–53, 2022.
- [3] Manuel Kauers and Jakob Moosbauer. Flip graphs for matrix multiplication. in preparation.
- [4] Julian D. Laderman. A noncommutative algorithm for multiplying  $3 \times 3$  matrices using 23 multiplications. *Bulletin of the American Mathematical Society*, 82(1):126–128, 1976.
- [5] François Le Gall. Powers of tensors and fast matrix multiplication. In *Proceedings of the 39th International Symposium on Symbolic and Algebraic Computation, ISSAC'14*, pages 296–303. ACM, 2014.
- [6] Alexandre Sedoglavic. Yet another catalogue of fast matrix multiplication algorithms. <https://fmm.univ-lille.fr/>, 2022. Accessed: 2022-10-08.
- [7] Alexandre Sedoglavic and Alexey V. Smirnov. The tensor rank of  $5 \times 5$  matrices multiplication is bounded by 98 and its border rank by 89. In *Proceedings of ISSAC'21*, pages 345–351, 2021.
- [8] Volker Strassen. Gaussian elimination is not optimal. *Numerische Mathematik*, 13(4):354–356, 1969.

MANUEL KAUERS, INSTITUTE FOR ALGEBRA, J. KEPLER UNIVERSITY LINZ, AUSTRIA  
 Email address: [manuel.kauers@jku.at](mailto:manuel.kauers@jku.at)

JAKOB MOOSBAUER, INSTITUTE FOR ALGEBRA, J. KEPLER UNIVERSITY LINZ, AUSTRIA  
 Email address: [jakob.moosbauer@jku.at](mailto:jakob.moosbauer@jku.at)