

Creative Telescoping



Manuel Kauers · Institute for Algebra · JKU

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- Such a recurrence is obtained from a **k-free recurrence** for $f(n, k)$.
- More precisely, we construct an annihilating operator for $f(n, k)$ of the form

$$P(n, S_n) - \Delta_k Q(n, k, \Delta_k, S_n)$$

where

- P is **nonzero** and **free of k, S_k, Δ_k** .
- Q may be zero and may involve any variables or operators.

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where

- P is **nonzero** and **free of k, S_k, Δ_k** .
- Q may be zero and may involve any variables or operators.
- Then $P(n, S_n)$ is an annihilating operator for $S(n)$.

$$P(n, S_n) - \Delta_k Q(n, k, \Delta_k, S_n)$$

“Telescopier”



$$P(n, S_n) - \Delta_k Q(n, k, \Delta_k, S_n)$$

“Telescoper”



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“Certificate”

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- A hypergeometric term $f(n, k)$ is said to **telescope** (w.r.t. k) if there is a hypergeometric term $g(n, k)$ such that

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- If we only have P , it is not obvious whether $P \cdot f$ telescopes, i.e., whether P is really a telescoper.
- If we have both P and Q , then checking $(P - \Delta_k Q) \cdot f \stackrel{?}{=} 0$ is easy. Thus Q **certifies** that P is a telescoper.

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$$((n+1)S_n - (4n+2)) \cdot \binom{n}{k}^2 = \Delta_k \cdot \frac{k^2(2k-3n-3)}{(n+1)^2} \binom{n+1}{k}^2$$

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The operator $\frac{k^2(2k-3n-3)}{(n+1)^2} S_n$ is a **certificate** for the telescoper.

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This recurrence together with the initial values $W_0 = \frac{\pi}{2}$ and $W_1 = 1$ determines the whole sequence.

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It maps this function to one that can be explicitly integrated.

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- $\sum_{n=0}^N (2 \binom{2n}{n} x^n - (1-4x) \binom{2n}{n} (x^n)') = (N+1) \binom{2N+2}{N+1} x^{N+1}$

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This differential equation together with the initial value $S(0) = 1$ implies the claimed identity.

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Therefore:

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This differential equation together with the initial values $K(0) = \frac{\pi}{2}$, $K'(0) = \frac{\pi}{8}$ uniquely describes $K(t)$.

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A Proof that Euler Missed ...

Apéry's Proof of the Irrationality of $\zeta(3)$

An Informal Report

Alfred van der Poorten



1. *Journées Arithmétiques de Marseille-Luminy, June 1978*

The board of programme changes informed us that K. Apéry (Cans) would speak Thursday, 14.00 "Sur l'irrationalité de $\zeta(3)$ ". Though there had been earlier rumours of his choosing a proof, suspicion was proved. The lecture tended to strengthen this view to rank disbelief. Those who listened casually, or who were afflicted with being non-Francoisophones, appeared to have only a glimpse of solidly assertions.

Exercice

Prove the following amazing claim:

① For all $\epsilon_1, \epsilon_2, \dots$

$$\sum_{k=1}^{\infty} \frac{\epsilon_1 \epsilon_2 \dots \epsilon_k}{(k + \epsilon_1) \dots (k + \epsilon_k)} > \frac{1}{2}$$

$$\textcircled{2} \zeta(3) = \sum_{n=1}^{\infty} \frac{1}{n^3} = \frac{5}{2} \sum_{n=1}^{\infty} \frac{(-1)^{n-1}}{n^3 \binom{2n}{n}}$$
 (1)

③ Consider the recursion:

$$a^2 b_n = (n-1)^2 b_{n-2} + (3n^2 - 3n^2 + 2n - 5) b_{n-1}, \quad n \geq 2. \quad (2)$$

Let $\{b_n\}$ be the sequence defined by $b_0 = 1, b_1 = 5$, and $b_n = a_n$ for all n ; then the b_n all are integers! Let $\{a_n\}$ be the sequence defined by $a_0 = 0, a_1 = 6$, and $a_n = a_n$ for all n ; then the a_n are rational numbers with denominator dividing $1, 2, 3, \dots, n!$ since $\{1, 2, \dots, n\}$ is the bin (Pascal's rule) of $\binom{2n}{n}$.

Exercice (continued)

$$\textcircled{4} \zeta(2) = \sum_{n=1}^{\infty} \frac{1}{n^2} = \frac{5}{6}$$

⑤ Consider the recursion:

$$n^2 b_n = (n-1)^2 b_{n-2} + (11n^2 - 11n + 7) b_{n-1}, \quad n \geq 2. \quad (3)$$

Let $\{b_n\}$ be the sequence defined by $b_0^2 = 1, b_1^2 = 5$ and the recursion; then the b_n^2 all are integers! Let $\{a_n^2\}$ be the sequence defined by $a_0^2 = 0, a_1^2 = 5$ and the recursion; then the a_n^2 are rational numbers with denominator dividing $1, 2, 3, \dots, n!$.

⑥ $a_n b_n \zeta(2) = \zeta(2) = a^2 b_n$; indeed the convergence is so fast as to imply that for all integers p, q with q sufficiently large relative to $\epsilon > 0$

$$|a^2 - \frac{p}{q}| > \frac{\epsilon}{q^{3/2}}, \quad q^2 = 11.85078 \dots$$

I heard with some incredulity that, for one, Houdouin (Bordeaux, now Grenoble) believed that these claims might be valid. Very much intrigued, I joined Hendrik Lenstra (Amsterdam) and Cohen in an evening's discussion in which Cohen explained and demonstrated most of the details of the proof. We came away convinced that Professor Apéry had indeed found a quite intricate and significant demonstration of the irrationality of $\zeta(3)$. But we remained unable to prove a critical step.

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① Consider the recurrence:

$$\zeta(3) - \sum_{n=1}^{\infty} \frac{1}{n^3} = 0$$
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$$\sum_{k=1}^{\infty} \frac{a_{k n_1} \dots a_{k n_2} \dots}{(k + n_1) \dots (k + n_2) \dots} = \frac{1}{\zeta(n_1) \dots \zeta(n_2)}$$

$$\zeta(3) - \sum_{n=1}^{\infty} \frac{1}{n^3} = \frac{5}{2} = \sum_{n=1}^{\infty} \frac{(1-1/n^3)^{-1}}{n^3} \quad (2)$$

② Consider the recurrence:

$$n^3 a_n - (n-1)^3 a_{n-1} = (34n^2 - 31n^2 + 23n - 5)a_{n-1}$$
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Let $\{b_n\}$ be the sequence defined by $b_0 = 1, b_1 = 5$, and $b_n = a_n$ for all n ; then the b_n all are integers! Let $\{c_n\}$ be the sequence defined by $c_0 = 0, c_1 = 6$, and $c_n = a_n$ for all n ; then the c_n are rational numbers with denominator dividing $2! \cdot 3! \dots n!$ since $\{1, 2, \dots, n\}$ is the bin (Pascal) rows of \dots

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$$|a^n - \frac{p}{q}| > \frac{1}{q^{3/2}}, \quad q^n > 11.85078 \dots$$

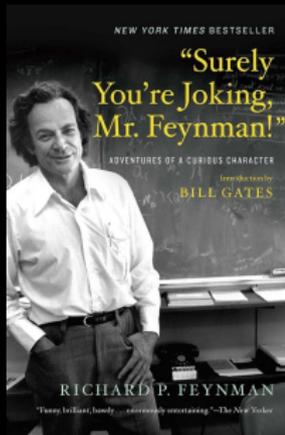
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① For all $n \geq 1$, $a_n, b_n \dots$

$$\frac{1}{k+1} \sum_{i=1}^k \frac{a_i \dots a_{i-1}}{(i+a_1) \dots (i+a_k)} = \frac{1}{a_{k+1}}$$

$$\textcircled{2} \zeta(3) = \sum_{n=1}^{\infty} \frac{1}{n^3} = \frac{5}{2} \sum_{n=1}^{\infty} \frac{(-1)^{n-1}}{n^3 \binom{2n}{n}}$$
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⑤ Consider the recursion:

$$n^3 a_n = (n-1)^3 a_{n-1} + (11n^2 - 11n + 7) a_{n-2}, \quad n \geq 2. \quad (6)$$

Let $\{b_n\}$ be the sequence defined by $b_0^2 = 1, b_1^2 = 5$ and the recursion; then the b_n^2 are integers! Let $\{c_n\}$ be the sequence defined by $c_0^2 = 0, c_1^2 = 5$ and the recursion; then the c_n^2 are rational numbers with denominator dividing $(1 \cdot 2 \cdot \dots \cdot n)^2$.

⑥ $a_n b_n = (n-2)2^n = n^2 6^n$, indeed the convergence is so fast as to imply that for all integers p, q with q sufficiently large relative to $c > 0$:

$$|a^n - \frac{c}{n}| > \frac{1}{q^{3/2}}, \quad q^n = 11.85078 \dots$$

I heard with some incredulity that, for one, Houdouin (Bordeaux, now Grenoble) believed that these claims might well be valid. Very much intrigued, I joined Hendrik Lenstra (Amsterdam) and Cohen in an evening discussion in which Cohen explained and demonstrated most of the details of the proof. We came away convinced that Professor Apéry had indeed found a quite intricate and significant demonstration of the irrationality of $\zeta(3)$. But we remained unable to prove a critical step.

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In the present segment, let's focus on the differential case, and rational functions.

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- **Find:** operators $P(x, D_x) \neq 0$ and $Q(x, y, D_x, D_y)$ such that

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f is integrable in $C(y)$ if and only if $h = 0$.

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Idea: find a $\mathbb{C}(x)$ -linear relation among h_0, h_1, h_2, \dots

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$$D_x^2 \cdot f = D_y \cdot \frac{-162x^3y^5+\dots-8y^2-16y-8}{x^2(27x+4)^2(xy^3+y+1)^2} + \frac{6(54x-27xy-y-1)}{x(27x+4)^2(xy^3+y+1)}$$

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 & + x(27x + 4) \frac{6(54x-27xy-y-1)}{x(27x+4)^2(xy^3+y+1)} = 0
 \end{aligned}$$

Therefore,

$$\begin{aligned}
 & (6 + 2(27x + 1)D_x + x(27x + 4)D_x^2) \cdot f \\
 & = D_y \cdot \frac{2(3x^2y^6+3xy^6-21xy^4-21xy^3-y^4-y^3+3y^2+6y+3)}{(xy^3+y+1)^3} + 0
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OUTPUT: a telescoper for f .

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Exercise. In step 2, we can use $D_x \cdot h_{r-1}$ instead of $D_x^r \cdot f$. Why is this better?

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For every $a_0, a_1, \dots, a_r \in C(x)$ (free of y), we have

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Then, with $s = (r - 1) \deg_y q + \deg_y p + 1$,

$$D_y \cdot g = \frac{\text{[redacted]}}{q^{r+1}} \quad \leftarrow \deg_y \leq r \deg_y q + \deg_y p$$

For undetermined a_0, \dots, a_r and b_0, \dots, b_s , enforce

$$(a_0 + a_1 D_x + \dots + a_r D_x^r) \cdot f \stackrel{!}{=} D_y g$$

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Comparing coefficients with respect to y leads to a linear system with

- $(r + 1) + (s + 1)$ variables
- $1 + r \deg_y q + \deg_y p$ equations

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Exercise. Does every nonzero solution give rise to a nonzero telescoper?

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The y -derivative of $g = \frac{b_0 + b_1 y + b_2 y^2 + b_3 y^3 + b_4 y^4}{(xy^3 + y + 1)^2}$ is

$$\frac{-2b_4 xy^6 - 3b_3 xy^5 + (2b_4 - 4b_2 x)y^4 + (b_3 - 5b_1 x + 4b_4)y^3 + (3b_3 - 6b_0 x)y^2 + (2b_2 - b_1)y - 2b_0 + b_1}{(xy^3 + y + 1)^3}$$

Example. $f = \frac{1}{xy^3 + y + 1}$

$$\begin{aligned} & (a_0x^2 - a_1x + 2a_2 + 2b_4x)y^6 \\ & \quad + 3b_3xy^5 \\ & + (2a_0x - a_1 + 4b_2x - 2b_4)y^4 \\ & + (2a_0x - a_1 + 5b_1x - b_3 - 4b_4)y^3 \\ & \quad + (a_0 + 6b_0x - 3b_3)y^2 \\ & \quad + (2a_0 + b_1 - 2b_2)y \\ & \quad + a_0 + 2b_0 - b_1 = 0 \end{aligned}$$

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$$\begin{pmatrix} x^2 & -x & 2 & 0 & 0 & 0 & 0 & 2x \\ 0 & 0 & 0 & 0 & 0 & 0 & 3x & 0 \\ 2x & -1 & 0 & 0 & 0 & 4x & 0 & -2 \\ 2x & -1 & 0 & 0 & 5x & 0 & -1 & -4 \\ 1 & 0 & 0 & 6x & 0 & 0 & -3 & 0 \\ 2 & 0 & 0 & 0 & 1 & -2 & 0 & 0 \\ 1 & 0 & 0 & 2 & -1 & 0 & 0 & 0 \end{pmatrix} \begin{pmatrix} a_0 \\ a_1 \\ a_2 \\ b_0 \\ b_1 \\ b_2 \\ b_3 \\ b_4 \end{pmatrix} = 0$$

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$$\begin{pmatrix} a_0 \\ a_1 \\ a_2 \\ b_0 \\ b_1 \\ b_2 \\ b_3 \\ b_4 \end{pmatrix} \in \left\langle \begin{pmatrix} 6x \\ 2x(27x + 1) \\ x^2(27x + 4) \\ -1 \\ 2(3x - 1) \\ 9x - 1 \\ 0 \\ -3x(x + 1) \end{pmatrix} \right\rangle$$

Example. $f = \frac{1}{xy^3 + y + 1}$

$$\begin{aligned} & (6x + 2x(27x + 1)D_x + x^2(27x + 4)D_x^2) \cdot f \\ &= D_y \cdot \frac{-1 + 2(3x - 1)y + (9x - 1)y^2 - 3x(x + 1)y^4}{(xy^3 + y + 1)^2} \end{aligned}$$

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OUTPUT: a telescoper and a certificate for f

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- 6 for a solution vector $(a_0, \dots, a_r, b_0, \dots, b_s)$ with at least one nonzero a_i , return P and g

Summary.

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$$P(\mathbf{n}, S_{\mathbf{n}}) - \Delta_{\mathbf{k}} Q(\mathbf{n}, \mathbf{y}, \Delta_{\mathbf{k}}, S_{\mathbf{n}})$$

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$$\underbrace{P(\mathbf{n}, S_n)}_{\text{telescoper (nonzero!)}} - \Delta_k \underbrace{Q(\mathbf{n}, \mathbf{y}, \Delta_k, S_n)}_{\text{certificate}}$$

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- **Creative telescoping** is the search for a telescoper (with or without a corresponding certificate).
- **Sister Celine's algorithm** is a creative telescoping algorithm for hypergeometric terms.
- For rational functions in the differential case, we have two creative telescoping algorithms:

Reduction-based telescoping

Hermite reduction
+ small linear system

Apagodu-Zeilberger algorithm

large linear system